## Multiplicative Data Perturbation for Privacy Preserving Data Mining

#### Kun Liu

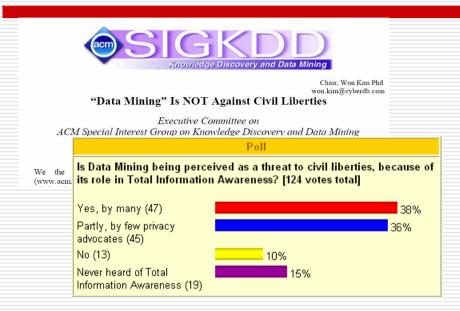
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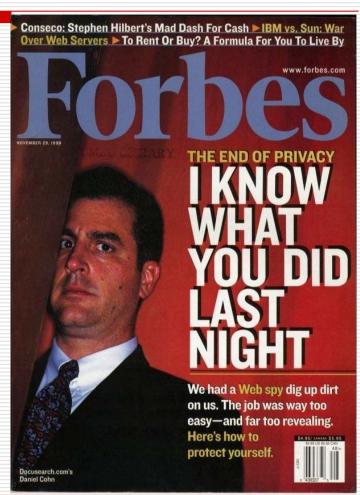


## **Growing Privacy Concerns**



"Detailed information on an individual's credit, health, and financial status, on characteristic purchasing patterns, and on other personal preferences is routinely recorded and analyzed by a variety of governmental and commercial organizations."

- M. J. Cronin, "e-Privacy?" Hoover Digest, 2000.



# Privacy-Preserving Data Mining

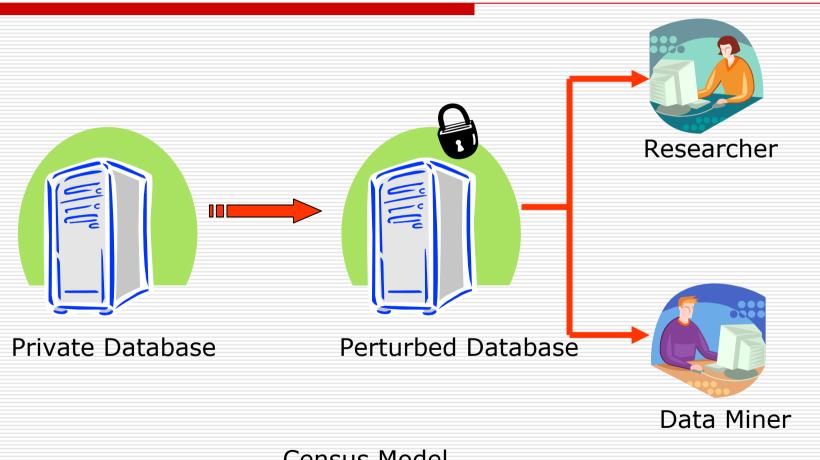
- "the best (and perhaps only) way to overcome the 'limitations' of data mining techniques is to do more research in data mining, including areas like data security and privacy-preserving data mining, which are actually active and growing research areas."
  - SIGKDD Executive Committee, "'Data Mining' Is NOT Against Civil Liberties," 2003.
- Privacy-preserving data mining is "the study of how to produce valid mining models and patterns without disclosing private information."
  - F. Giannotti and F. Bonchi, "Privacy Preserving Data Mining," KDUbiq Summer School, 2006.

## Privacy-Preserving Data Mining

- Data Perturbation
  - Hiding private data while mining patterns
- Secure Multi-Party Computation
  - Building a model over multi-party distributed databases without knowing others' inputs
- Knowledge Hiding
  - Hiding sensitive rules/patterns
- Privacy-aware Knowledge Sharing
  - Do the data mining results themselves violate privacy?

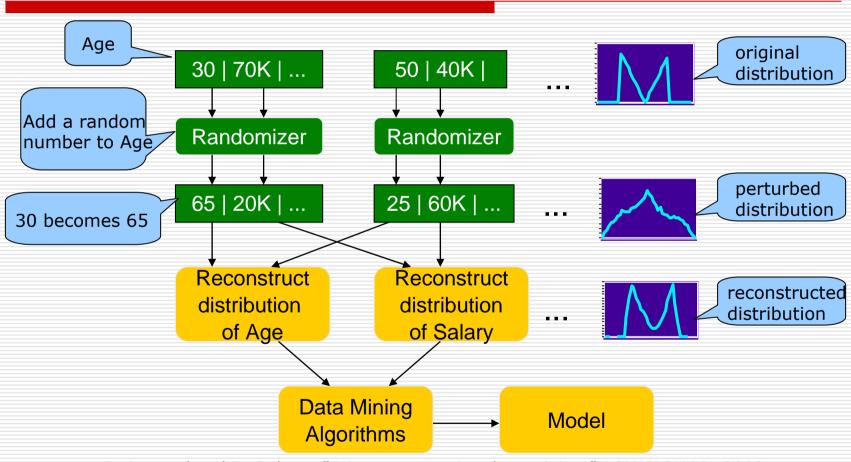
[more]

## Data Perturbation



Census Model

## Additive Data Perturbation



## Additive vs. Multiplicative Noise

- □Additive perturbation is not safe.
  - ☐ "in many cases, the original data can be accurately estimated from the perturbed data using a spectral filter that exploits some theoretical properties of random matrices"
  - Kargupta et al., "On the Privacy Preserving Properties of Random Data Perturbation Techniques," ICDM, 2003.
  - ☐ Related work: [Huang05], [Guo06], etc.
- ■How about multiplicative noise?
  - Has not been carefully studied.
  - Topic of this dissertation.

## **Primary Contributions**

- We examined the effectiveness of exact Euclidean distance preserving data perturbation, and developed three attack techniques.
  - K. Liu, C. Giannella, and H. Kargupta, "An attacker's view of distance preserving maps for privacy preserving data mining," 10th European Conference on Principles and Practice of Knowledge Discovery in Databases (PKDD'06), 2006.
- We proposed a random projection-based approximate distance preserving perturbation as a possible remedy, and analyzed its privacy issues.
  - K. Liu, H. Kargupta, and J. Ryan, "Random projection-based multiplicative perturbation for privacy preserving distributed data mining," *IEEE Transactions on Knowledge and Data Engineering (TKDE)*, 18(1), 2006.

## Roadmap

- □ Traditional Multiplicative Noise
- Distance Preserving Data Perturbation
  - Fundamental Properties
  - Known Input-Output Attack
  - Know Sample Attack
  - Independent Signal Attack
- Random Projection-based Perturbation
  - Fundamental Properties
  - Bayes Privacy Model
  - Attacks Revisit
- Conclusion and Future Work

# Traditional Multiplicative Noise

#### Private Database X

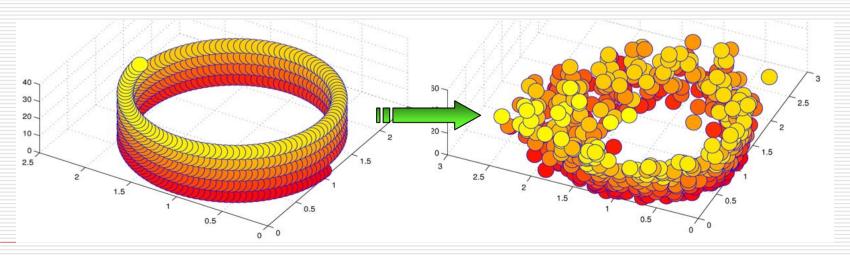
ID	1001	1002
Wages	98,563	83,821
Rent	1,889	1,324
Tax	2,899	2,578

#### Perturbed Database Y

ID	1001	1002
Wages	116,166	85,396
Rent	1,878	1,381
Tax	2,964	2,135

2,899 \* 1.0224 = 2,964

 $y_{ij} = x_{ij} \times r_{ij}$ , where  $x_{ij}$  is the private data,  $r_{ij} \sim N(1, \sigma)$  [Kim03].



## Traditional Multiplicative Noise

- Mechanism
  - Each data element/cell is randomized independently by multiplying a random number.
- Pros
  - Summary statistics (e.g., mean, variance) can be estimated from the perturbed data.
  - Effective if data disseminator only wants minor perturbation
  - Popular in the statistics community.
- Cons
  - Equivalent to additive perturbation after a logarithmic operation. Vulnerable to attacks designed for additive noise.
  - Not preserving Euclidean distance; not suitable for many data mining tasks.

## Roadmap

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- □ Distance Preserving Data Perturbation
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## Distance Preserving Perturbation

Dist. preserving perturbation

$$T: \mathbb{R}^n \to \mathbb{R}^n \text{ if } \forall x, y \in \mathbb{R}^n, \|x - y\| = \|T(x) - T(y)\|$$

- Dist. preserving perturbation is equivalent to
  - $x \in \mathbb{R}^n \to Mx + v$ , for  $M \in \mathcal{O}_n$  and  $v \in \mathbb{R}^n$ ,
  - where  $O_n$  is the set of all  $n \times n$  orthogonal matrices.
- Dist. preserving perturbation with origin fixed
  - $x \in \mathbb{R}^n \to Mx$ , where  $M \in \mathcal{O}_n \to \mathcal{O}$  Orthogonal Transformation

**Our Focus** 

## Distance Preserving Perturbation

ID	1001	1002
Wages	-26,326	-22,613
Rent	-94,502	-80,324
Tax	10,151 8,432	

	-0.2507	0.4556	-0.8542
=	-0.9653	-0.0514	0.2559
	0.0726	0.8887	0.4527

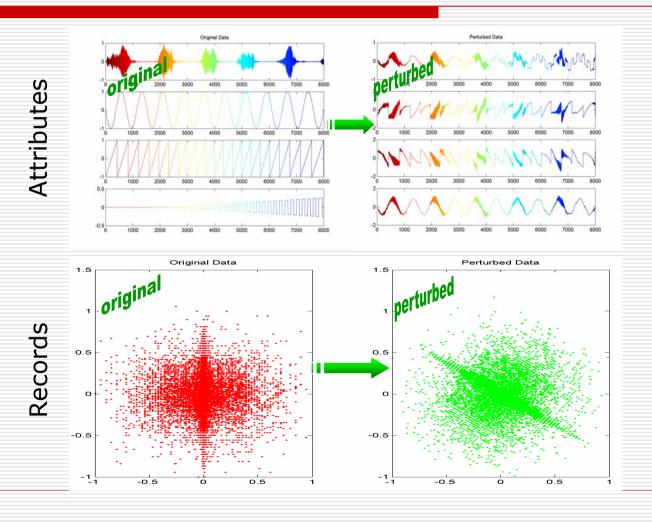
Μ

ID	1001	1002
Wages	98,563	83,821
Rent	1,889	1,324
Tax	2,899	2,578

Υ

- $\square$  Perturbation Model Y = MX
  - X: original private data with each column a record
  - Y: perturbed data
  - M: orthogonal perturbation matrix
- Perturbed data produces exactly the same data mining results
  - Clustering [Oliveira04], Classification [Chen05]
  - Other related: [Mukherjee06], etc.

## Distance Preserving Perturbation



### Is Dist. Preserving Perturbation Secure?

- Known Input-Output Attack: attacker knows some collection of linearly independent private data records and their corresponding perturbed version.
- Known Sample Attack: attacker has a collection of independent data samples from the same distribution the original data was drawn.
- Independent Signals Attack: all data attributes are non-Gaussian and statistically independent

## Privacy Breach

ε-Privacy Breach

For any  $\varepsilon > 0$ , we say that an  $\varepsilon$ -privacy breach occurs if

$$\|\hat{x} - x_i\| \le \|x_i\| \varepsilon$$

where  $\hat{x}$  is the attacker's estimate of  $x_i$ , the  $i^{th}$  data tuple in X.

Probability of ε-Privacy Breach

$$\rho(x_i, \varepsilon) = \text{Prob}\{\|\hat{x} - x_i\| \le \|x_i\| \varepsilon\}$$

the probability that an  $\varepsilon$ -privacy breach occurs.

## Known Input-Output Attack

$$[Y_{n\times k} \quad Y_{n\times (m-k)}] = M_{n\times n} [X_{n\times k} \quad X_{n\times (m-k)}]$$
KNOWN

- $\square$  Assumption (can be relaxed): rank( $X_{nxk}$ )=k
- $\square$  If k=n:
  - $M = Y_{n \times k} X^{-1}_{n \times k}, \ X_{n \times (m-k)} = M^{T} Y_{n \times (m-k)}$
  - Probability of privacy breach  $\rho(x_i, \varepsilon) = 1$  for  $\varepsilon = 0$  and any i.
  - The attacker has a perfect recovery of the private data.
- ☐ If k<n, what is going to happen?</p>

## Known Input-Output Attack

$$[Y_{n\times k} \quad Y_{n\times (m-k)}] = M_{n\times n} [X_{n\times k} \quad X_{n\times (m-k)}]$$

 $\square$  If k<n, any matrix  $\hat{M}$  in the set

$$\Omega = \{ \hat{M} \in \mathcal{O}_n : \hat{M}X_{n \times k} = Y_{n \times k} \}$$

can be the original perturbation matrix  $M_{n \times n}$ , where is  $O_n$  is the set of all nxn orthogonal matrices.

The attacker chooses one uniformly from  $\Omega$  as an estimation of  $M_{n\times n}$ , uses that to recover other private data, and computes the probability of privacy breach.

## Known Input-Output Attack

#### Probability of Privacy Breach

$$\rho(x_{i}, \varepsilon) = \operatorname{Prob}\{ \| \hat{x} - x_{i} \| \leq \| x_{i} \| \varepsilon \}$$

$$= \operatorname{Prob}\{ \| \hat{M}Mx_{i} - x_{i} \| \leq \| x_{i} \| \varepsilon \}$$

$$= \begin{cases} \frac{1}{\pi} 2 \arcsin\left(\frac{\| x_{i} \| \varepsilon}{2d(x_{i}, X_{n \times k})}\right) & \text{if } \| x_{i} \| \varepsilon < 2d(x_{i}, X_{n \times k}) ; \\ 1 & \text{otherwise.} \end{cases}$$

where  $d(x_i, X_{n \times k})$  is the distance of  $x_i$  from the column space of  $X_{n \times k}$ , and  $\hat{M}$  is uniformly chosen from  $\Omega = \{\hat{M} \in O_n : MX_{n \times k} = Y_{n \times k}\}$ .

#### Known Input-Output Attack Example

Private Data X: X₁->Y₁ KNOWN	X <sub>1</sub> /25.0000 75.0000	X <sub>2</sub> 30.0000 90.0000	X <sub>3</sub> 45.0000 105.0000	)—→UNKNOWN
Perturbed Data Y:	Y <sub>1</sub> -42.0198 66.9652	Y <sub>2</sub> -50.4237 80.3582	Y <sub>3</sub> -68.5443 91.3875	

- The distance of  $X_2$  from the column space of  $X_1$  is 0, therefore  $\rho(x_2, \varepsilon) = 1$  for any  $\varepsilon$ .
- The distance of  $X_3$  from the column space of  $X_1$  is 9.4868, therefore  $\rho(x_3, \varepsilon) = \frac{1}{\pi} 2 \arcsin\left(\frac{\|x_3\| \varepsilon}{2 \times 9.4868}\right)$ , e.g.  $\rho(x_3, 0.01) = 3.84\%$ .

## Known Sample Attack

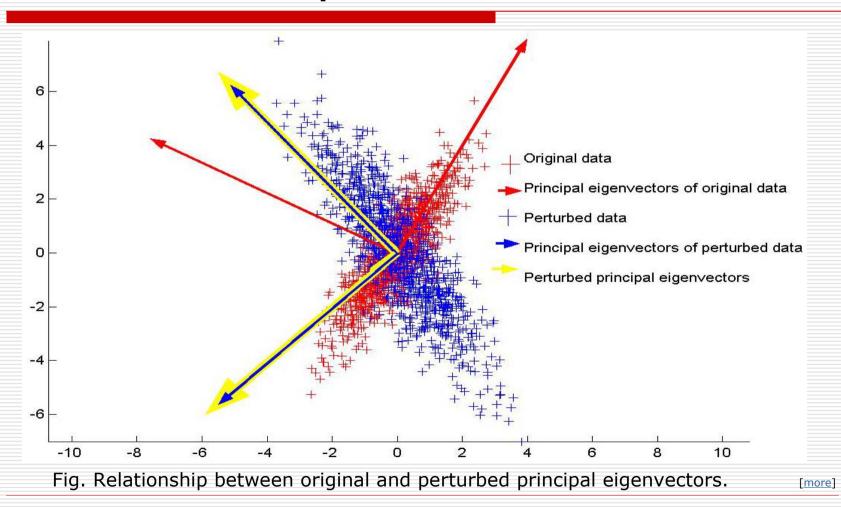
#### Assumptions

- Each data record arose as an independent sample from some unknown distribution
- The attacker has a collection of samples independently chosen from the same distribution
- The covariance of the distribution has all distinct eigenvalues (holds true in most practical situations [Jolliffe02]).

#### Attack Technique

Exploring the relationship between the principal eigenvectors of the original data and the principal eigenvectors of the perturbed data.

## Known Sample Attack



## Known Sample Attack Experiments

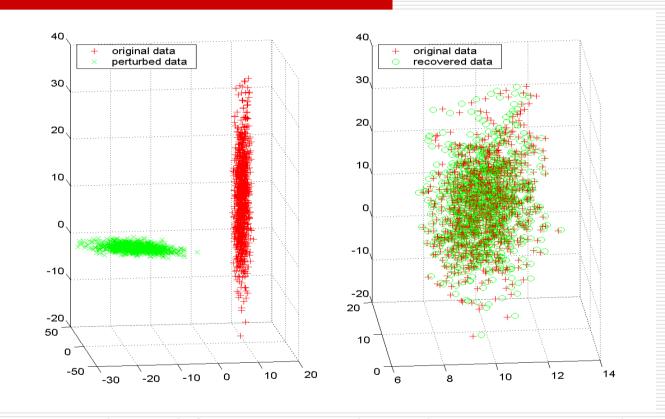


Fig. Known sample attack for 3D Gaussian data with 10,000 private tuples. The attacker has 2% samples from the same distribution. The average relative error of the recovered data is 0.0265 (2.65%).

## Known Sample Attack Experiments

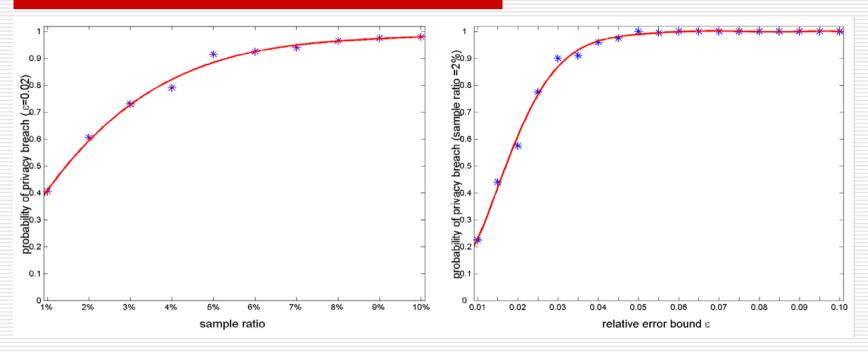


Fig. Probability of privacy breach w.r.t. attacker's sample size. The relative error bound  $\epsilon$  is fixed to be 0.02. (3D Gaussian data with 10,000 private tuples.)

Fig. Probability of privacy breach w.r.t. the relative error bound  $\,\varepsilon$  . The sample ratio is fixed to be 2%. (3D Gaussian data with 10,000 private tuples.)

more

## Independent Signals Attack

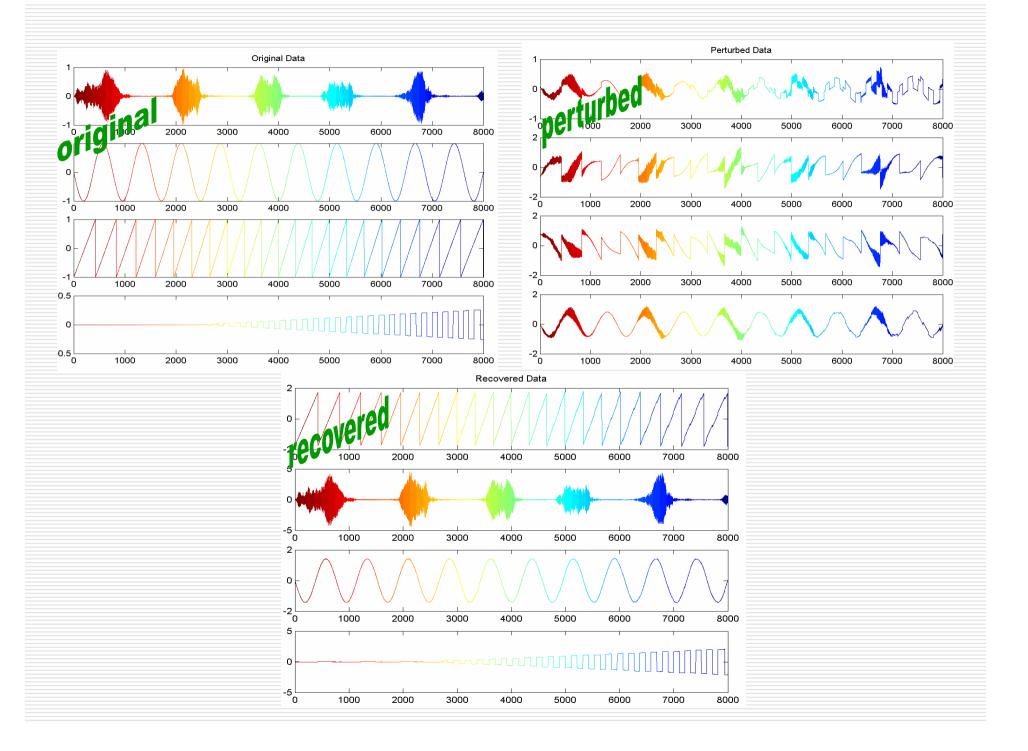
□ Basic Independent Component Analysis Model

- □ Objective: recover the original signals X from only the observed mixtures Y.
- Requirements
  - Source signals are statistically independent
  - All signals must be non-Gaussian with exception of one
  - k ≥ n
  - Matrix M must be of full column rank

more

## Independent Signals Attack Experiments





#### Distance Preserving Perturbation Summary

- Mechanism
  - Whole data set is perturbed by multiplying an orthogonal matrix.
- □ Pros
  - Perturbed data preserves Euclidean distance.
  - Many data mining algorithms can be applied to the perturbed data and produce exactly the same results as if applied to the original data.
- Cons
  - Vulnerable to known input-output attack
  - Vulnerable to known sample attack
  - Vulnerable to independent signals attack

## Roadmap

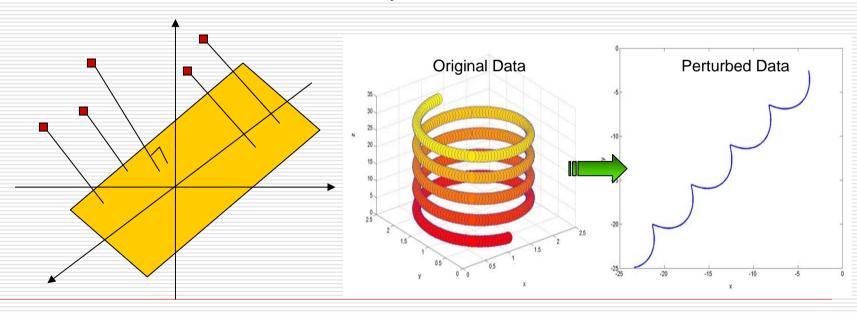
- Traditional Multiplicative Noise
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## Random Projection

#### Basic Model

$$u = \frac{1}{\sqrt{k}\sigma_r} R_{k \times m} x_{m \times 1}, \text{ and } v = \frac{1}{\sqrt{k}\sigma_r} R_{k \times m} y_{m \times 1},$$

where k < m and  $r_{ij}$  is i.i.d.  $\sim N(0, \sigma_r)$ .



## Random Projection

Preserving Inner Product
$$E[u^{T}v - x^{T}y] = 0 \text{ and } Var[u^{T}v - x^{T}y] = \frac{1}{k} \left(\sum_{i} x_{i}^{2} \sum_{i} y_{i}^{2} + \left(\sum_{i} x_{i} y_{i}\right)^{2}\right).$$

→The distortion produced by random projection is zero on the average, and its variance is inversely proportional to k, dimension of new space.

[more]

#### Preserving Euclidean Distance

$$\Pr\{(1-\eta) \| x - y \|^2 \le \| u - v \|^2 \le (1+\eta) \| x - y \|^2\} = \int_{k(1-\eta)}^{k(1+\eta)} f(t;k) dt, \ \eta > 0$$
 where  $f(t;k)$  is the p.d.f. of chi-square distribution with  $k$  degrees of freedom.

 $\rightarrow$ The probability that relative error is bounded with in (1±  $\eta$ ) increases proportionally with k. more ]

## Bayes Privacy Model

- Primitives
  - Let x be the private data and y the perturbed one.
  - Attacker's Prior Belief: f(x)
  - Attacker's Additional Background Knowledge: θ
  - Attacker's Posteriori Belief: f(x | y, θ )
- Information Non-Disclosure Principle
  - The perturbed data should provide the attacker with little additional information beyond the attacker's prior belief and other background knowledge.
- Example
  - ( $\rho_1$ ,  $\rho_2$ ,)-privacy [Evfimevski03] happens when  $f(x) < \rho_1$  and  $f(x | y, \theta) > \rho_2$  OR  $f(x) > 1 \rho_1$  and  $f(x | y, \theta) < 1 \rho_2$

## Maximum a Posteriori Probability (MAP) Estimate

- $\square$  ( $\rho_1$ ,  $\rho_2$ , )-privacy works only for discrete data. It assumes statistically independent inputs and outputs, and requires transition probability explicitly defined. Not appropriate for multiplicative perturbation.
- We propose a maximum a posteriori probability (MAP) estimate-based approach
  - $1. \hat{x}_{MAP} = \arg\max_{x} f(x \mid y, \theta)$
  - **2.**  $\hat{x}_{MAP}$  is compared with  $x_i$  to see whether any extra information is disclosed, e.g.,  $\|\hat{x}_{MAP} x_i\| \le \|x_i\| \varepsilon$ .

# Why Maximum a Posteriori Probability (MAP) Estimate

- It is closely related to maximum a posteriori probability hypothesis testing.
  [more]
- ☐ It considers both prior and posterior belief. In the absence of a priori knowledge, MAP estimate becomes maximum likelihood estimate (MLE).
- ☐ It often produces estimates with errors that are not much higher than the minimum mean square error.
- It is relatively easy to derive the conditional p.d.f. in the multiplicative data perturbation scenario.

## Maximum a Posteriori Probability (MAP) Estimate

- Assumption I: Attackers' best knowledge of f(x) is it is uniformly distributed over an multi-dimensional interval.
- Assumption II: Attacker has no other background knowledge, i.e.,  $\theta = \emptyset$ .
- MAP Estimate:  $\hat{x}_{MAP} = \arg \max_{x} f(x \mid y, \theta)$   $= \arg \max_{x} \frac{k^{1/2}}{(2\pi x^{T} x)^{k/2}} \exp\left(-\frac{ky^{T} y}{2x^{T} x}\right),$ where  $x \in \mathbb{R}^{n}$  and  $y \in \mathbb{R}^{k}$ .
- □ Solution: Any  $\hat{x}$  in the interval that satisfies  $\hat{x}^T \hat{x} = y^T y$ .
- Conclusion: MAP does not offer attacker more info than what has been implied by properties of random projection itself.

# Privacy / Accuracy Control

Random Projection

$$u = \frac{1}{\sqrt{k}\sigma_r} R_{k \times m} x_{m \times 1}, \text{ and } v = \frac{1}{\sqrt{k}\sigma_r} R_{k \times m} y_{m \times 1},$$

where k < m and  $r_{ii}$  is i.i.d.  $\sim N(0, \sigma_r)$ .

Accuracy

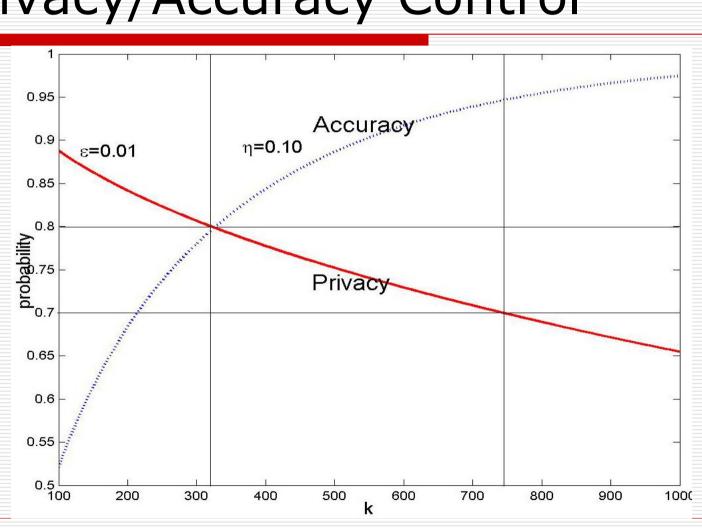
$$\Pr\{(1-\eta)\|x-y\|^2 \le \|u-v\|^2 \le (1+\eta)\|x-y\|^2\} = \int_{k(1-\eta)}^{k(1+\eta)} f(t;k)dt, \ \eta > 0.$$

□ ¬ ε-Privacy Breach

$$\Pr\{\|\hat{x}_{MAP} - x\| > \|x\| \varepsilon\} = \int_{-\infty}^{k(1-\varepsilon)^2} f(t;k)dt + \int_{k(1+\varepsilon)^2}^{+\infty} f(t;k)dt.$$

Here f(t;k) is the p.d.f. of chi-square distribution with k degrees of freedom.

# Privacy/Accuracy Control



#### Roadmap

- Traditional Multiplicative Noise
- Distance Preserving Data Perturbation
  - Fundamental Properties
  - Known Input-Output Attack
  - Know Sample Attack
  - Independent Signal Attack
- □ Random Projection-based Perturbation
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# Independent Signals Attack

$$Y_{k \times m} = \frac{1}{\sqrt{k}\sigma_r} R_{k \times n} X_{n \times m}$$

- When k<n, at most (k-1) source signals can be separated out [Cao96].</p>
- With probability one, linear ICA can't separate out any of the original signals for any  $(k \times n)$   $(k \le n/2, n \ge 2)$  random matrix with i.i.d. entries chosen from continuous distribution [Liu06a].

#### Independent Signals Attack Experiments

original





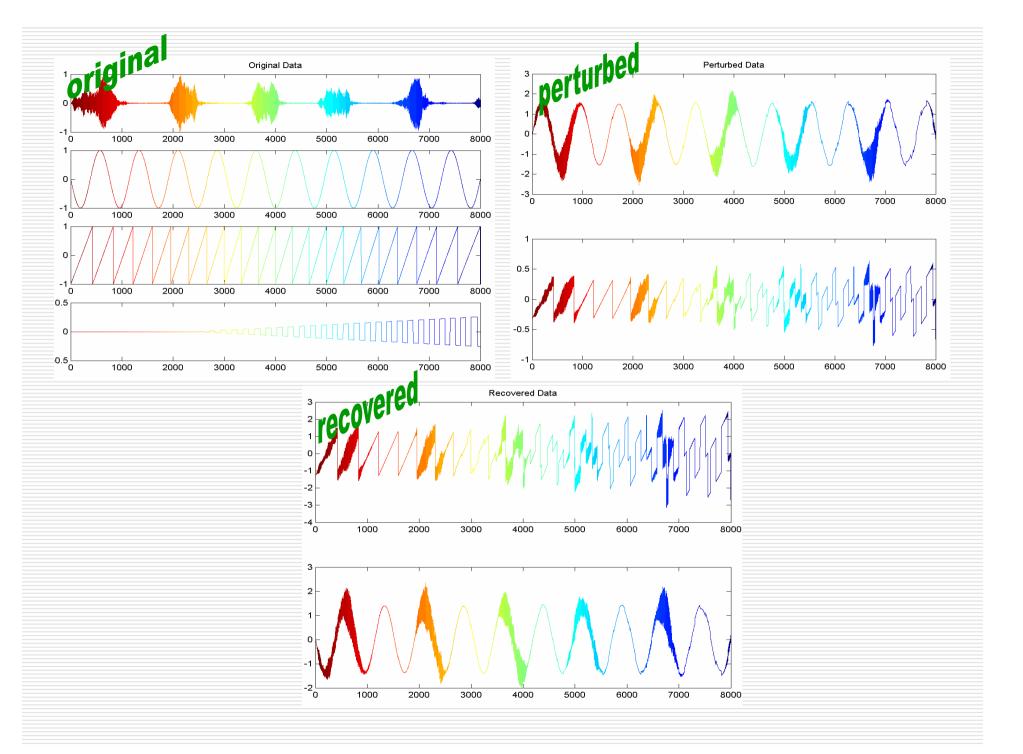




perturbed

recovered





# Known Sample Attack

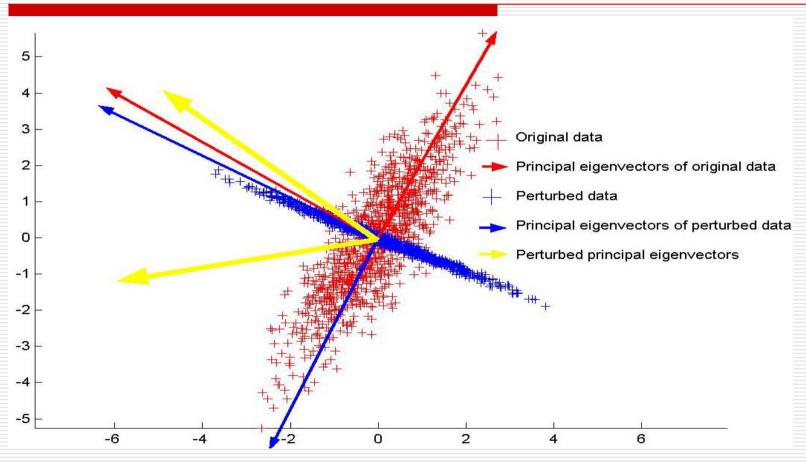


Fig. Relationship between original and perturbed principal eigenvectors.

# Known Input-Output Attack

$$[Y_{k\times p} \quad Y_{k\times (m-p)}] = \frac{1}{\sqrt{k}\sigma_r} R_{k\times n} [X_{n\times p}] \quad X_{n\times (m-p)}]$$
KNOWN

- If p=n and  $rank(X_{nxp}) = p$ , R can be recovered, but still it is an under-determined system of linear equations.
- MAP estimate shows that relative error decreases as known input-output pairs increases; relative error increases as k decreases.

# Random Projection-based Perturbation Summary

- Mechanism
  - Data is projected to a lower dimensional random space.
- Pros
  - From the perspective of MAP estimate, random projection does not disclose more information than what have been implied by the distance preservation properties.
  - It offers better privacy protection than orthogonal transformation-based distance preserving perturbation.
- Cons
  - Perturbed data approximately preserves Euclidean distance, therefore little loss in accuracy.

#### Conclusions

- Traditional Multiplicative Data Perturbation
- Distance Preserving Data Perturbation
  - Known Input-Output Attack (linear algebra, statistics)
  - Known Sample Attack (PCA)
  - Independent Signals Attack (ICA)
- □ Random projection-based Data Perturbation
  - Accuracy Analysis
  - Bayes Privacy Model
  - Maximum a Posteriori Probability (MAP) Estimate
  - Privacy/Accuracy Control
  - Attack Analysis
- ☐ Privacy Issues Are Intrinsically Complex
  - Need joined efforts from researchers, engineers, sociologists, legal experts, policy makers...

#### Future Work

- A game theoretic framework for large scale distributed privacy preserving data mining
  - Distributed and ubiquitous computing becomes popular
  - Some participants cooperative and honest, some malicious
  - Computation in such environment is more like a game
  - Necessary to develop a game theoretic framework
- Combination of cryptographic techniques and perturbation techniques
  - Cryptographic techniques offers strong privacy guarantee, but with high communication and computation cost
  - Perturbation provides statistically weaker privacy protection, but more efficient
  - Would be ideal to combine them to achieve both efficiency and privacy

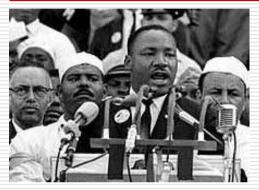
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- [Huang05] Z. Huang, W. Du, and B. Chen, "Deriving private information from randomized data," in Proceedings of the 2005 ACM SIGMOD Conference (SIGMOD'05), Baltimore, MD, June 2005, pp. 37-48.
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#### I Have a Dream



"I have a dream that one day this nation will rise up and live out the true meaning of its creed: 'We hold these truths to be self-evident, that all men are created equal."

- Martin Luther King, Jr., 1963.

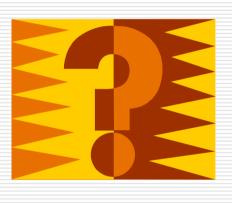


"I have a dream that one day I will get a Ph.D. degree."

- Kun Liu, when he was a kid.

# Thank you and Questions

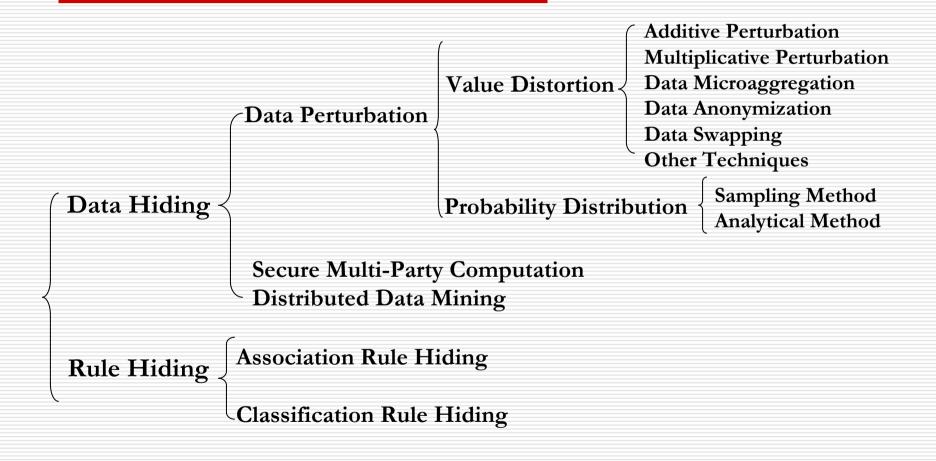






# Backup Slides

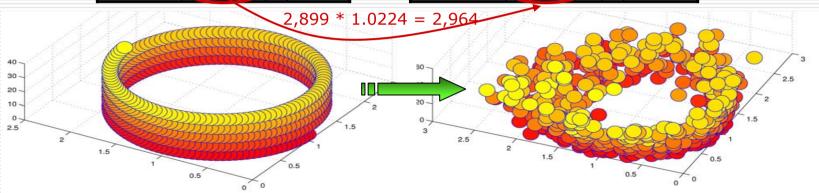
#### Overview of PPDM backup



# Traditional Multiplicative Noise backup

ID	1001	1002
Wages	98,563	83,821
Rent	1,889	1,324
Tax	2 899	2 578

ID	1001	1002
Wages	116,166	85,396
Rent	1,878	1,381
Tax	2,964	2,135



#### Properties:

- $y_{ij} = x_{ij} \times r_{ij}$ , where  $x_{ij}$  is the private data,  $r_{ij} \sim N(1, \sigma)$  [Kim03].
- Each data element randomized independently.
- Original Mean and variance can be estimated from perturbed data.
- Equivalent to additive perturbation after a logarithmic operation.
- Not preserve Euclidean distance.

# Knowledge Hiding backup

- What is disclosed?
  - the data (modified somehow)
- What is hidden?
  - some "sensitive" knowledge (i.e. secret rules/patterns)
- □ How?
  - usually by means of data sanitization. The data which we are going to disclose is modified, in such a way that the sensitive knowledge can no longer be inferred, while the original database is modified as less as possible.

# Privacy-aware Knowledge Sharing backup

- What is disclosed?
  - the intentional knowledge (i.e., rules, patterns, models)
- What is hidden?
  - the source data
- The central question
  - Do the data mining results themselves violate privacy

#### Privacy-aware Knowledge Sharing backup

```
Age = 27, Zip = 15254, Christian->American
(sup count = 758, confidence = 99.8\%)
```

```
Age = 27, Zip = 15254->American
(sup count = 1518, confidence = 99.9\%)
```

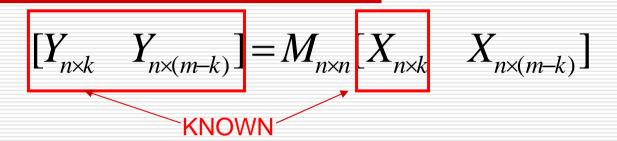
```
sup_count (27, 15254, Christian) = 758/.998 = 759.5
sup_count(27, 15254, Christian, 7 American) = 759.5*0.002 = 1.519
```

```
\sup_{x \in \mathbb{R}} count(27, 15254) = 1518/0.999 = 1519.5
sup\_count(27, 15254, \neg American) = 1519.5 * 0.001 = 1.5195
```

```
Age = 27, Postcode = 45254, ¬ American->Christian
(sup count \approx 1.5, confidence \approx 100.0%)
```

This information refers to my France neighbor.... he is Christian!

# Known Input-Output Attack

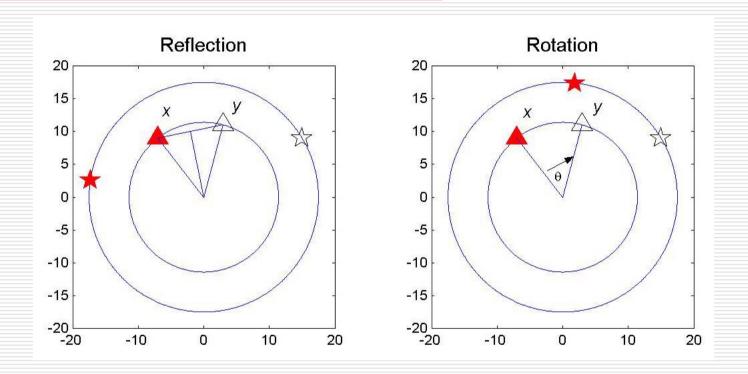


 $\square$  Closed-form Expression of  $\Omega$ 

$$\Omega = \{ \hat{M} \in \mathcal{O}_n : MX_{n \times k} = Y_{n \times k} \}$$
$$= \{ M (U_k U_k^T + U_{n-k} P U_{n-k}^T) : \forall P \in \mathcal{O}_{n-k} \},$$

where  $U_k$  is the orthonormal basis for the column space of  $X_{n \times k}$ ,  $U_{n-1}$  is the orthonormal basis for the orthogonal complement of the column space of  $X_{n \times k}$ .

# Known Input-Output Attack backup



Special case in 2D space: when k = 1 and n = 2. The attacker can't distinguish rotation and reflection.

# Known Input-Output Attack backup

- Properties of the Probability of Privacy Breach
  - Attacker can compute the probability of privacy breach for a given private record and a relative error bound  $\mathcal{E}$  .
  - The larger the  $\mathcal E$  , the higher the probability of privacy breach.
  - The closer the private record is to the column space of the known records, the higher the probability of privacy breach.
  - The distance  $d(x_i, X_{n \times k})$  can be computed from the perturbed data.

# Known Sample Attack

The principal eigenvectors of the original data have experienced the same distance preserving perturbation as the data itself.

```
Let Y = MX, we have Z_Y = MZ_XD,
where Z_{v} is the eigenvector matrix of the covariance of Y;
Z_{x} is the eigenvector matrix of the covariance of X;
and D is a diagonal matrix with each entry on the diagonal \pm 1.
```

- $Z_v$  can be computed from the perturbed data,  $Z_x$  can be estimated from the sample data. (See dissertation for choice of D, details omitted.)
- Attacker uses  $Z_x$ ,  $Z_y$  and D to recover M, and therefore X.

#### Known Sample Attack Experiments Lackup

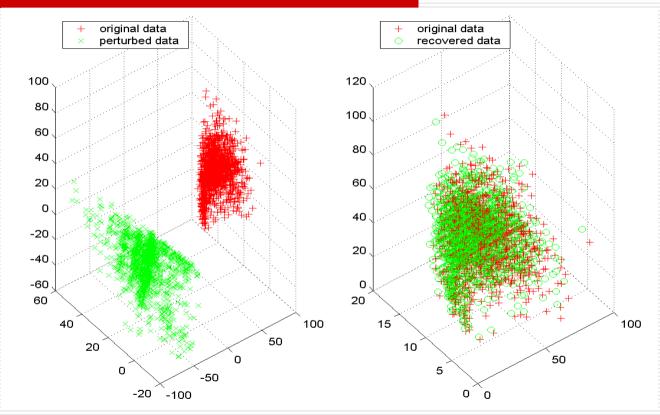


Fig. Known sample attack for Adult data with 32,561 private tuples. The attacker has 2% samples from the same distribution. The average relative error of the recovered data is 0.1081 (10.81%).

#### Known Sample Attack Experiments backup



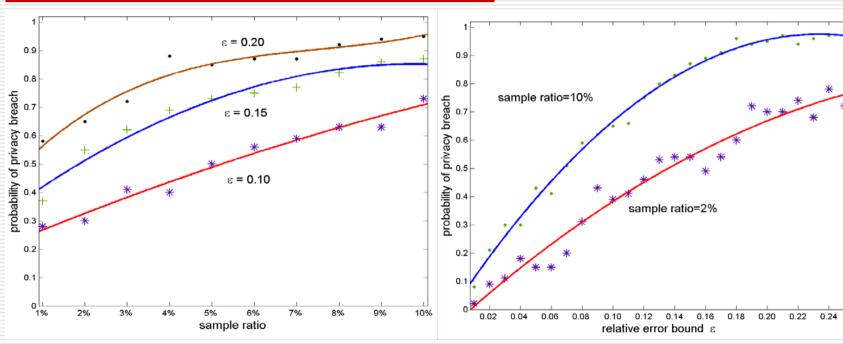


Fig. Probability of privacy breach w.r.t. attacker's sample size. The relative error bound  $\varepsilon$  changes from 0.10 to 0.20. (Adult data with 32,561 private tuples)

Fig. Probability of privacy breach w.r.t. the relative error bound  $\epsilon$  . The sample ratio is fixed to be 2% and 10%. (Adult data with 32,561 private tuples.)

#### Effectiveness of Known Sample Attack

- Covariance Estimation Quality
  - Larger sample size gives attacker better recovery
  - Robust covariance estimator helps to downweight the influence of outliers
- p.d.f. of the Data
  - The greater the difference between any pair of eigenvalues of the covariance, the higher the probability of privacy breach
- More details can be found in the dissertation.

#### Independent Component Analysis

- $\begin{bmatrix} m_{11} & m_{12} & \dots & m_{1n} \end{bmatrix} \begin{bmatrix} x_1(t_1) & x_1(t_2) & \dots & x_1(t_m) \end{bmatrix}$ Basic Model  $Y_{k \times m} = M_{k \times n} X_{n \times m} = \begin{vmatrix} m_{21} & m_{22} & \dots & m_{2n} \\ \dots & \dots & \dots & \dots \end{vmatrix} \begin{vmatrix} x_2(t_1) & x_2(t_2) & \dots & x_2(t_m) \\ \dots & \dots & \dots & \dots & x_3(t_m) \end{vmatrix}$  $m_{k1}$   $m_{k2}$  ...  $m_{kn} \mid x_n(t_1) x_n(t_2)$  ...  $x_n(t_m)$
- □ ICA Estimation
  - To find a matrix W such that WY = WMX = X
- Nongaussian is Independent
  - Central limit theory sum of random variables has a distribution closer to Gaussian than any of the original random variables.
  - ICA looks for a W that maximizes the nongaussianity of WY.
- Measures of Nongaussianity
  - Kurtosis:  $kurt(x) = E[x^4] 3E^2[x^2]$
  - Negentropy:  $J(x) = H(x_{gaussian}) H(x)$
  - Mutual information:  $I(x_1, x_2, ..., x_n) = \sum_{i=1}^{n} H(x_i) H(x)$

# Random Projection

#### Relative Errors in Computing the Inner Product of Two Attributes

k	Mean(%)	Var(%)	Min(%)	Max(%)
100(1%)	9.91	0.41	0.07	23.47
500(5%)	5.84	0.25	0.12	18.41
1000(10%)	2.94	0.05	0.03	7.53
2000(20%)	2.69	0.04	0.01	7.00
3000(30%)	1.81	0.03	0.27	6.32

#### Relative Errors in Computing the Euclidean Distance of the Two Attributes

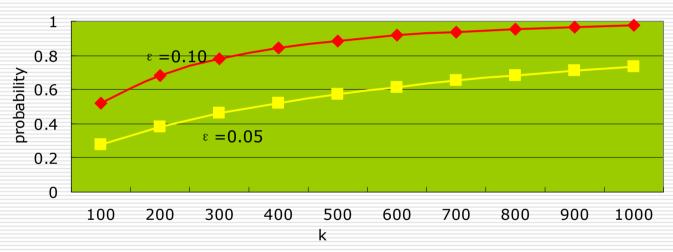
k	Mean(%)	Var(%)	Min(%)	Max(%)
100(1%)	10.44	0.67	1.51	32.58
500(5%)	4.97	0.29	0.23	18.32
1000(10%)	2.70	0.05	0.11	7.21
2000(20%)	2.59	0.03	0.31	6.90
3000(30%)	1.80	0.01	0.61	3.91

Adult data from UCI Repository. The first 10,000 elements of attributes fnlwgt and education-num.

# Random Projection

backup

 $\Pr\{(1-\eta) \| x - y \|^2 \le \| u - v \|^2 \le (1+\eta) \| x - y \|^2\} = \int_{k(1-\eta)}^{k(1+\eta)} f(t;k) dt, \ \eta > 0$  where f(t;k) is the p.d.f. of chi-square distribution with k degrees of freedom.



The probability of the accuracy of random projection w.r.t. k and  $\,\epsilon$ . Each entry of the random matrix is i.i.d., chosen from a Gaussian distribution with mean zero and constant variance.

The probability that relative error is bounded with in  $(1 \pm \eta)$  increases proportionally with k.

#### Maximum a Posteriori (MAP) Test backup

Given a binary hypothesis testing experiment with outcome s, the following rule leads to the lowest possible value of P<sub>FRROR</sub>:

```
s \in A_0 if Prob\{H_0 \mid s\} \ge Prob\{H_1 \mid s\}; s \in A_1 otherwise.
```

- Here  $P_{ERROR} = \text{Prob}\{A_1 \mid H_0\} \text{ Prob}\{H_0\} + \text{Prob}\{A_0 \mid H_1\} \text{ Prob}\{H_1\}.$
- The test design divides S into two sets,  $A_0$  and  $A_1 = A_0^c$ . If the outcome s is in  $A_0$ , the conclusion is accept  $H_0$ . Otherwise, the conclusion is accept  $H_1$ .

#### MAP Known Input-Output Attack backup

$$[Y_{k\times p} \quad Y_{k\times (m-p)}] = \frac{1}{\sqrt{k}\sigma_r} R_{k\times n} [X_{n\times p} \quad X_{n\times (m-p)}]$$
KNOWN

- Assumption I: Attackers' best knowledge of f(X) is it is uniform.
- Assumption II: Attacker has no other background knowledge, *i.e.*,  $\theta = \emptyset$ .

$$\hat{x}_{MAP} = \arg\max_{x} f(\mathbf{x} = x \mid \frac{1}{\sqrt{k}} \mathbf{R} \mathbf{x} = y, \frac{1}{\sqrt{k}} \mathbf{R} X_{p} = Y_{p}),$$

$$= \arg\max_{x} (2\pi)^{-\frac{1}{2}k(p+1)} \det\left(\frac{1}{k} \overline{X}^{T} \overline{X}\right)^{-\frac{1}{2}k} etr\left(-\frac{1}{2} \overline{Y} (\frac{1}{k} \overline{X}^{T} \overline{X})^{-1} \overline{Y}^{T}\right),$$

$$where \ r_{ij} \sim N(0,1), \ \overline{X} = [x \ X_{p}], \ \overline{Y} = [y \ Y_{p}], \ \overline{X} \ has full \ column \ rank.$$