# CMSC 471 Fall 2012 

## Class \#9

## Thurs 9/27/12 Game Theory

Kevin Winner, winnerk1@umbc.edu

## Today's class

- Game playing
- Nim
- Stochastic games
- Game Theory


## Game Playing \#2

## Still chapter 5

Some material adopted from notes by Charles R. Dyer, University of Wisconsin-Madison

## Example: Nim

- In Nim, there are a certain number of objects (coins, sticks, etc.) on the table -- we'll play 7-coin Nim
- Each player in turn has to pick up either one or two objects
- Whoever picks up the last object loses


## Games of chance

- Backgammon is a two-player game with uncertainty.
-Players roll dice to determine what moves to make.
-White has just rolled 5 and 6 and has four legal moves:

$$
\begin{array}{r}
\cdot 5-10,5-11 \\
\cdot 5-11,19-24 \\
\cdot 5-10,10-16 \\
\cdot 5-11,11-16
\end{array}
$$

- Such games are good for exploring decision making in adversarial problems involving skill and luck.



## Game trees with chance nodes

- Chance nodes (shown as circles) represent random events
- For a random event with N outcomes, each chance node has N distinct children; a probability is associated with each
- (For 2 dice, there are 21 distinct outcomes)
- Use minimax to compute values for MAX and MIN nodes
- Use expected values for chance nodes
- For chance nodes over a max node, as in C :
$\operatorname{expectimax}(\mathrm{C})=\sum_{\mathrm{i}}\left(\mathrm{P}\left(\mathrm{d}_{\mathrm{i}}\right) *\right.$ maxvalue $\left.(\mathrm{i})\right)$
- For chance nodes over a min node:
$\operatorname{expectimin}(\mathrm{C})=\sum_{\mathrm{i}}\left(\mathrm{P}\left(\mathrm{d}_{\mathrm{i}}\right) *\right.$ minvalue $\left.(\mathrm{i})\right)$
MAX

DICE

MIN

DICE

MAX

TERMINAL


## Example: Oopsy-Nim

- Starts out like Nim
- Each player in turn has to pick up either one or two objects
- Sometimes (with probability 0.25 ), when you try to pick up two objects, you drop them both
- Picking up a single object always works
- Whoever picks up the last object loses
- Question: Why can't we draw the entire game tree?


## Game Theory

## Not actually in your textbook

## Game Theory

- Reasoning about multi-agent interactions
- Games are episodic (although sometimes we'll track an interaction history)
- Agents select moves simultaneously and independently
- Describe the game as a table of each permutation of actions
- Based on our predictions of what the other agent(s) will do, what is the optimal strategy for us to play?


## Rock-Paper-Scissors

|  |  | Player 2 |  |  |
| :---: | :---: | :---: | :---: | :---: |
|  |  | R | P | S |
| Player 1 | P | $1,-1$ | $-1,1$ | $1,-1$ |
|  | R | 0,0 | 0,0 | $-1,1$ |
|  | S | $-1,1$ | $1,-1$ | 0,0 |

## Prisoner's Dilemma

- From Wikipedia:
-"Two men are arrested, but the police do not have enough information for a conviction. The police separate the two men, and offer both the same deal: if one testifies against his partner (defects/betrays), and the other remains silent (cooperates with/assists his partner), the betrayer goes free and the one that remains silent gets a one year sentence. If both remain silent, both are sentenced to only one month in jail on a minor charge. If each 'rats out' the other, each receives a three-month sentence. Each prisoner must choose either to betray or remain silent; the decision of each is kept secret from his partner. What should they do?"


## Prisoner's Dilemma



## Prisoner's Dilemma



## What is the best strategy?



## Pareto Optimality

- $S$ is a Pareto-optimal solution iff
$-\forall S^{\prime}\left(\exists_{x} U_{x}\left(S^{\prime}\right)>U_{x}(S) \rightarrow \exists y U_{y}\left(S^{\prime}\right)<U_{y}(S)\right)$
- i.e., if $X$ is better off in $S^{\prime}$, then some $Y$ must be worse off
- Social welfare, or global utility, is the sum of all agents' utility
- If S maximizes social welfare, it is also Pareto-optimal (but not vice versa)



## Stability

- If an agent can always maximize its utility with a particular strategy (regardless of other agents' behavior) then that strategy is dominant
- A set of agent strategies is in Nash equilibrium if each agent's strategy Si is locally optimal, given the other agents' strategies
- No agent has an incentive to change strategies
- Hence this set of strategies is locally stable


## Game Analysis

- Which solution(s) maximizes social welfare?
- Which solution(s) are Paretooptimal?
- Which solution(s) are Nash equilibriums?
- What is the dominant strategy(ies)?

| 3,2 | 1,1 | 4,4 | 0,5 | 1,1 |
| :---: | :---: | :---: | :---: | :---: |
| 1,7 | 1,2 | 0,0 | 7,8 | 1,5 |
| 2,1 | 2,1 | 0,0 | 2,9 | 2,1 |
| 6,6 | 0,0 | 3,3 | 5,6 | 2,3 |
| 4,0 | 3,10 | 5,3 | 2,10 | 6,1 |

## Iterated Prisoner's Dilemma (IPD)

- Play Prisoner's Dilemma many times with the same opponent


## Iterated Prisoner's Dilemma (IPD)

- Play Prisoner's Dilemma many times with the same opponent



## IPD Strategies

- Always Cooperate
- Always Defect
- Optimistic defect
- Tit-for-tat
- Tit-for-two-tats
- Tit-for-tat with forgiveness
- Master-slave
- And lots of more complicated strategies


## Other games we can play

- Ultimatum Game
- Two players are splitting $\$ 100$. One player offers a split to the other. If the split is accepted, both players receive payout equal to the offer. Otherwise, neither player receives anything.
- Guess $2 / 3$ of the average
- Everyone in the class submits a number from 0-100. The objective is to guess closest to $2 / 3$ of the average of all the guesses.
- Stag hunt

|  | Stag | Hare |
| :---: | :---: | :---: |
| Stag | 2,2 | 0,1 |
| Hare | 1,0 | 1,1 |

