# CMSC 471 Fall 2012 

## Class \#8

# Tue 9/25/12 Game Playing 

Kevin Winner, winnerk1@umbc.edu

# Game Playing 

## Chapter 5

Some material adopted from notes by Charles R. Dyer, University of Wisconsin-Madison

## Today's class

- Homework 2
- Game playing
- State of the art and resources
- Framework
- Game trees
- Minimax
- Alpha-beta pruning
- Stochastic games
- Homework 1
- Homework 2 design


## Why study games?

- Offer an opportunity to study interesting environments which are still easy to understand
-Multiagent, adversarial
- Stochastic
- Partially observable
- Clear criteria for success
- Games often define very large search spaces
-Chinook (checkers) has a dictionary of 39 trillion end-game states
- Fun


## State of the art

- How good are computer game players?
- Chess:
- Deep Blue beat Gary Kasparov in 1997
- Checkers: Chinook (an AI program with a very large endgame database) has solved checkers
- Go: Computer players have finally reached tournamentlevel play
- Backgammon: TD-Gammon trained on itself to reach the ability to beat top-level players
- Poker: After chess, probably the most studied game for AI players


## Typical case

- 2-person game
- Players alternate moves
- Zero-sum: one player's loss is the other's gain
- Perfect information: both players have access to complete information about the state of the game. No information is hidden from either player.
- Deterministic: no chance (e.g., using dice) involved
- Examples: Tic-Tac-Toe, Checkers, Chess, Go, Nim, Othello
- Not: Bridge, Solitaire, Backgammon, ...


## How to play a game

- A way to play such a game is to:
- Consider all the legal moves you can make
- Compute the new position resulting from each move
- Evaluate each resulting position and determine which is best
- Make that move
- Wait for your opponent to move and repeat
- Key problems are:
- Representing the "board"
- Generating all legal next boards
- Evaluating a position


## Evaluation function

- Evaluation function or static evaluator is used to evaluate the "goodness" of a game position.
- Contrast with heuristic search where the evaluation function was a non-negative estimate of the cost from the start node to a goal and passing through the given node
- The zero-sum assumption allows us to use a single evaluation function to describe the goodness of a board with respect to both players.
- $\mathbf{f ( n )} \gg 0$ 0: position $n$ good for me and bad for you
$-\mathbf{f}(\mathbf{n}) \ll \mathbf{0}$ : position n bad for me and good for you
- $\mathbf{f}(\mathbf{n})$ near 0: position $n$ is a neutral position
$-\mathbf{f}(\mathbf{n})=+$ infinity: win for me
$-\mathbf{f}(\mathbf{n})=$-infinity: win for you


## Evaluation function examples

- Example of an evaluation function for Tic-Tac-Toe:
$f(n)=[\#$ of 3-lengths open for me] - [\# of 3-lengths open for you]
where a 3-length is a complete row, column, or diagonal
- Alan Turing's function for chess
$-\mathbf{f}(\mathbf{n})=\mathbf{w}(\mathbf{n}) / \mathbf{b}(\mathbf{n})$ where $\mathrm{w}(\mathrm{n})=$ sum of the point value of white's pieces and $b(n)=$ sum of black's
- Most evaluation functions are specified as a weighted sum of position features:

$$
\mathrm{f}(\mathrm{n})=\mathrm{w}_{1} * \operatorname{feat}_{1}(\mathrm{n})+\mathrm{w}_{2} * \operatorname{feat}_{2}(\mathrm{n})+\ldots+\mathrm{w}_{\mathrm{n}} * \operatorname{feat}_{\mathrm{k}}(\mathrm{n})
$$

- Example features for chess are piece count, piece placement, squares controlled, etc.
- Deep Blue had over 8000 features in its evaluation function


## Gд!

- Problem spaces for typical games are ${ }_{\text {м쑝ㅇ }}$ represented as trees
- Root node represents the current board configuration; player must decide the best single move to make next
- Static evaluator function rates a board position. $\mathrm{f}($ board $)=$ real number with $\mathrm{f}>0$ "white" (me), $\mathrm{f}<0$ for black (you) ${ }_{\text {uma }}$

- Arcs represent the possible legal moves for a player
- If it is my turn to move, then the root is labeled a "MAX" node; otherwise it is labeled a "MIN" node, indicating my opponent's turn.
- Each level of the tree has nodes that are all MAX or all MIN; nodes at level i are of the opposite kind from those at level i+1


## Minimax procedure

- Create start node as a MAX node with current board configuration
- Expand nodes down to some depth (a.k.a. ply) of lookahead in the game
- Apply the evaluation function at each of the leaf nodes
- "Back up" values for each of the non-leaf nodes until a value is computed for the root node
- At MIN nodes, the backed-up value is the minimum of the values associated with its children.
- At MAX nodes, the backed-up value is the maximum of the values associated with its children.
- Pick the operator associated with the child node whose backed-up value determined the value at the root


## Minimax Algorithm



MAX
MIN

## Minimax Algorithm



MAX
MIN

## Minimax Algorithm



MAX
MIN

## Minimax Algorithm



## Minimax Algorithm



## Minimax Tree

MAX


MIN

## Minimax Tree

MAX

MIN


## Minimax Tree



## Minimax Tree



## Minimax Tree



## Minimax Tree



## Minimax Tree



## Minimax Tree



## Example: Nim

- In Nim, there are a certain number of objects (coins, sticks, etc.) on the table -- we'll play 7-coin Nim
- Each player in turn has to pick up either one or two objects
- Whoever picks up the last object loses


## Example: Nim

- In Nim, there are a certain number of objects (coins, sticks, etc.) on the table -- we'll play 7-coin Nim
- Each player in turn has to pick up either one or two objects
- Whoever picks up the last object loses


## Example: Nim

- In Nim, there are a certain number of objects (coins, sticks, etc.) on the table -- we'll play 7-coin Nim
- Each player in turn has to pick up either one or two objects
- Whoever picks up the last object loses


## Example: Nim

- In Nim, there are a certain number of objects (coins, sticks, etc.) on the table -- we'll play 7-coin Nim
- Each player in turn has to pick up either one or two objects
- Whoever picks up the last object loses


## Example: Nim

- In Nim, there are a certain number of objects (coins, sticks, etc.) on the table -- we'll play 7-coin Nim
- Each player in turn has to pick up either one or two objects
- Whoever picks up the last object loses


## Example: Nim

- In Nim, there are a certain number of objects (coins, sticks, etc.) on the table -- we'll play 7-coin Nim
- Each player in turn has to pick up either one or two objects
- Whoever picks up the last object loses


## Example: Nim

- In Nim, there are a certain number of objects (coins, sticks, etc.) on the table -- we'll play 7-coin Nim
- Each player in turn has to pick up either one or two objects
- Whoever picks up the last object loses


## Example: Nim

- In Nim, there are a certain number of objects (coins, sticks, etc.) on the table -- we'll play 7-coin Nim
- Each player in turn has to pick up either one or two objects
- Whoever picks up the last object loses


## Alpha-beta pruning

- We can improve on the performance of the minimax algorithm through alpha-beta pruning
- Basic idea: "If you have an idea that is surely bad, don't take the time to see how truly awful it is. " -- Pat Winston



## Alpha-beta pruning

- Traverse the search tree in depth-first order
- At each MAX node $n$, alpha(n) = maximum value found so far
- At each MIN node $n, \operatorname{beta}(\mathbf{n})=$ minimum value found so far
- Note: The alpha values start at -infinity and only increase, while beta values start at +infinity and only decrease.
- Beta cutoff: Given a MAX node $n$, cut off the search below $n$ (i.e., don't generate or examine any more of n's children) if alpha(n) $>=$ beta(i) for some MIN node ancestor i of $n$.
- Alpha cutoff: stop searching below MIN node n if beta(n) $<=$ alpha(i) for some MAX node ancestor $i$ of $n$.


## Alpha-beta example

MAX


MIN

## Alpha-beta example

MAX

MIN


## Alpha-beta example



## Alpha-beta example



## Alpha-beta example



## Alpha-beta example



## Alpha-beta example



## Alpha-beta example



## Alpha-beta example



## Alpha-beta example



## Alpha-beta example



## Alpha-beta example



## Alpha-beta example



## Alpha-beta algorithm

```
function MAX-VALUE (state, \alpha, \beta)
    ; ; 人 = best MAX so far; }\beta=\mathrm{ best MIN
if TERMINAL-TEST (state) then return UTILITY(state)
v := -\infty
for each s in SUCCESSORS (state) do
    v := MAX (v, MIN-VALUE ( s, \alpha, \beta))
    if v >= \beta then return v
    \alpha := MAX ( }\alpha,v
end
return v
function MIN-VALUE (state, 人, \beta)
if TERMINAL-TEST (state) then return UTILITY(state)
v := \infty
for each s in SUCCESSORS (state) do
    v := MIN (v, MAX-VALUE (s, \alpha, \beta))
    if v <= \alpha then return v
    \beta:= MIN ( }\beta,\textrm{v}
end
return v
```


## Games of chance

- Backgammon is a two-player game with uncertainty.
-Players roll dice to determine what moves to make.
-White has just rolled 5 and 6 and has four legal moves:

$$
\begin{array}{r}
\cdot 5-10,5-11 \\
\cdot 5-11,19-24 \\
\cdot 5-10,10-16 \\
\cdot 5-11,11-16
\end{array}
$$

- Such games are good for exploring decision making in adversarial problems involving skill and luck.



## Game trees with chance nodes

- Chance nodes (shown as circles) represent random events
- For a random event with N outcomes, each chance node has N distinct children; a probability is associated with each
- (For 2 dice, there are 21 distinct outcomes)
- Use minimax to compute values for MAX and MIN nodes
- Use expected values for chance nodes
- For chance nodes over a max node, as in C :
$\operatorname{expectimax}(\mathrm{C})=\sum_{\mathrm{i}}\left(\mathrm{P}\left(\mathrm{d}_{\mathrm{i}}\right) *\right.$ maxvalue $\left.(\mathrm{i})\right)$
- For chance nodes over a min node:
$\operatorname{expectimin}(\mathrm{C})=\sum_{\mathrm{i}}\left(\mathrm{P}\left(\mathrm{d}_{\mathrm{i}}\right) *\right.$ minvalue $\left.(\mathrm{i})\right)$
MAX

DICE

MIN

DICE

MAX

TERMINAL


## Meaning of the evaluation function



- Dealing with probabilities and expected values means we have to be careful about the "meaning" of values returned by the static evaluator.
- Note that a "relative-order preserving" change of the values would not change the decision of minimax, but could change the decision with chance nodes.
- Linear transformations are OK


## Example: Oopsy-Nim

- Starts out like Nim
- Each player in turn has to pick up either one or two objects
- Sometimes (with probability 0.25 ), when you try to pick up two objects, you drop them both
- Picking up a single object always works
- Whoever picks up the last object loses
- Question: Why can't we draw the entire game tree?


## Example: Oopsy-Nim

- Starts out like Nim
- Each player in turn has to pick up either one or two objects
- Sometimes (with probability 0.25 ), when you try to pick up two objects, you drop them both
- Picking up a single object always works
- Whoever picks up the last object loses
- Question: Why can't we draw the entire game tree?


## Example: Oopsy-Nim

- Starts out like Nim
- Each player in turn has to pick up either one or two objects
- Sometimes (with probability 0.25 ), when you try to pick up two objects, you drop them both
- Picking up a single object always works
- Whoever picks up the last object loses
- Question: Why can't we draw the entire game tree?


## Example: Oopsy-Nim

- Starts out like Nim
- Each player in turn has to pick up either one or two objects
- Sometimes (with probability 0.25 ), when you try to pick up two objects, you drop them both
- Picking up a single object always works
- Whoever picks up the last object loses
- Question: Why can't we draw the entire game tree?


## Example: Oopsy-Nim

- Starts out like Nim
- Each player in turn has to pick up either one or two objects
- Sometimes (with probability 0.25 ), when you try to pick up two objects, you drop them both
- Picking up a single object always works
- Whoever picks up the last object loses
- Question: Why can't we draw the entire game tree?

