
Duplicates are not allowed.

$$
\begin{aligned}
& \text { The elements in the BST must be comparable. }
\end{aligned}
$$

$10 / 12 / 2005$
BST Implementation
The SearchTree ADT

- A search tree is a binary search tree which stores
homogeneous elements with no duplicates.
- It is dynamic.
- The elements are ordered in the following ways
- inorder -- as dictated by operator<
- preorder, postorder, levelorder -- as dictated by the structure of
the tree
- Each BST maintains a simple object, known as
ITEM_NOT_FOUND, that is guaranteed to not be an
element of the tree. ITEM_NOT_FOUND is provided
to the constructor. (author's code)
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BinarySearchTree class


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BinarySearchTree class (cont) private:
BinaryNode
const Compa
const Compa
BinarySearchTree Implementation


Asymptotic performance is $O$ (height) in all cases
10/12/2005
Predecessor in BST
Predecessor of a node $v$ in a BST is the node that holds the
data value that immediately precedes the data at $v$ in order.

$$
\begin{aligned}
& \text { Finding predecessor } \\
& \text { - } \mathrm{v} \text { has a left subtree } \\
& \text { - then predecessor must be the largest value in the left subtree } \\
& \text { (the rightmost node in the left subtree) } \\
& -\mathrm{v} \text { does not have a left subtree } \\
& \text { - predecessor is the first node on path back to root that does not } \\
& \text { have } \mathrm{v} \text { in its left subtree }
\end{aligned}
$$

$$
\begin{aligned}
& \text { Successor of a node } \mathrm{v} \text { in a BST is the node that holds the data } \\
& \text { value that immediately follows the data at } \mathrm{v} \text { in order. } \\
& \text { Finding Successor } \\
& \text { - v has right subtree } \\
& \text { • successor is smallest value in right subtree } \\
& \text { (the leftmost node in the right subtree) } \\
& \text { - v does not have right subtree } \\
& \text { • successor is first node on path back to root that does not have } \mathrm{v} \\
& \text { in its right subtree }
\end{aligned}
$$

const
$+$
template <class Comparable>
void BinarySearchTree<Comparable>: :
remove (const Comparable\& x, BinaryNode<Comparable> *\& t) const
\&


The insert Operation

Tree Iterators
Could provide separate iterators for each desired order

- Iterator $\langle\mathbb{T}\rangle *$ GetInorderIterator ();
- Iterator $\langle\mathbb{T}\rangle *$ GetPreorderIterator ();
- Iterator $\langle\mathbb{T}\rangle *$ GetPostorderIterator ();
- Iterator $\langle\mathbb{T}\rangle *$ GetLevelorderIterator ();
14

Iterator<T> BinaryTree: :GetInorderIterator() \{
m_theList $=$ new ArrayList<T>;
FullListInorder (m_theList, getRoot ());
return m_theList->GetIterator ();


$$
\begin{aligned}
& \qquad \text { Tree Iterators (cont) } \\
& \text { Approach 2: store traversal in stack to mimic } \\
& \text { recursive traversal } \\
& \text { template <class } \mathrm{T}\rangle \\
& \text { class InorderIterator } \\
& \text { \{ private: } \\
& \text { Stack<* BNode<T\gg m_stack; }
\end{aligned}
$$

$$
\begin{aligned}
& \} ; \\
& 10 / 12 / 2005
\end{aligned}
$$


\}
$10 / 12 / 2005$
template <class $\mathrm{T}>$
T InOrderIterator<T
T InOrderIterator<T>: :Next ()

return top->element;
10/12/2005

