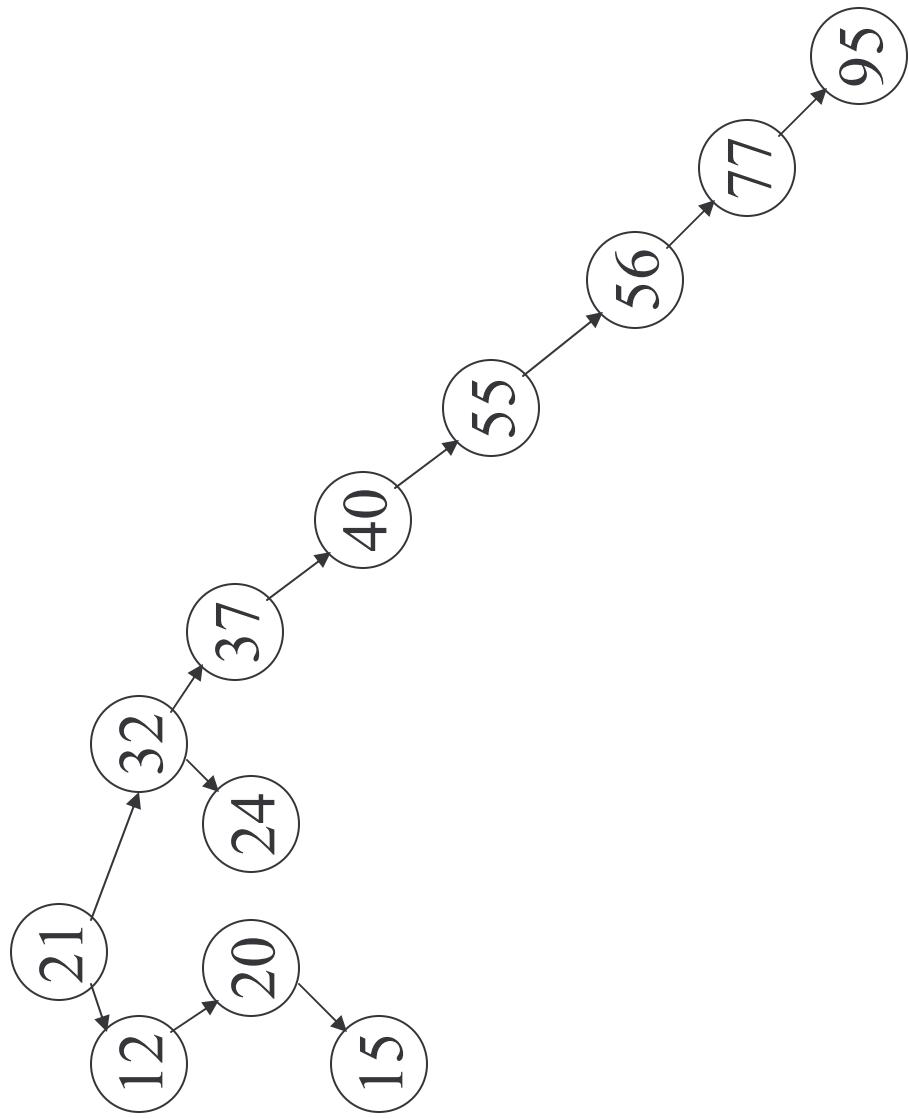


CMSC 341

Splay Trees

## Problems with BSTs

Because the shape of a BST is determined by the order that data is inserted, we run the risk of trees that are essentially lists



# Splay Trees

## Concept

- adjust tree in response to accesses to make common operations (insert, find, remove) efficient
- after access node is moved to root by *splaying*

## Performance

- amortized such that m operations take  $O(m \lg n)$  where n is the number of insertions (nodes in the tree)

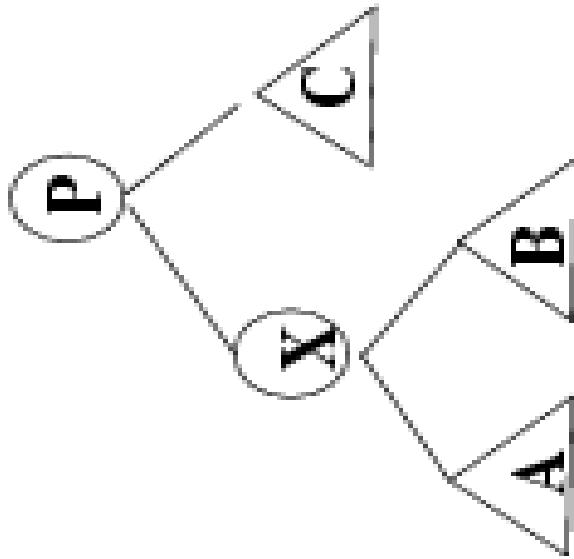
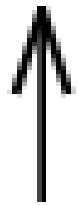
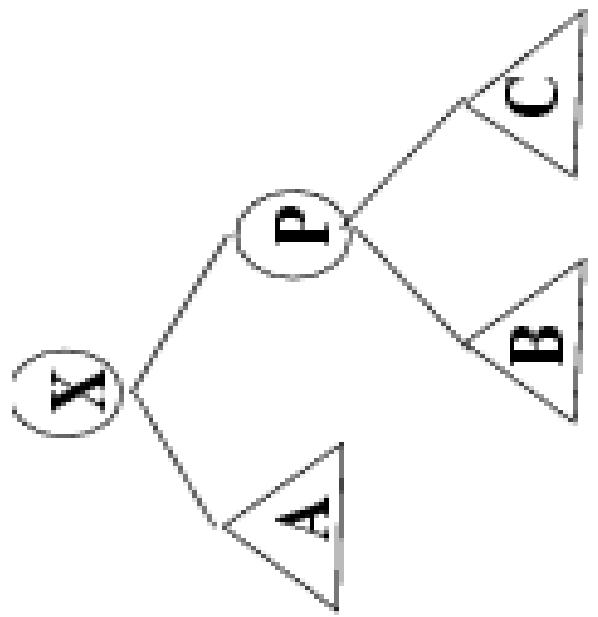
# Splay Operation

Traverse tree from node  $x$  to root, rotating along the way until  
 $x$  is the root

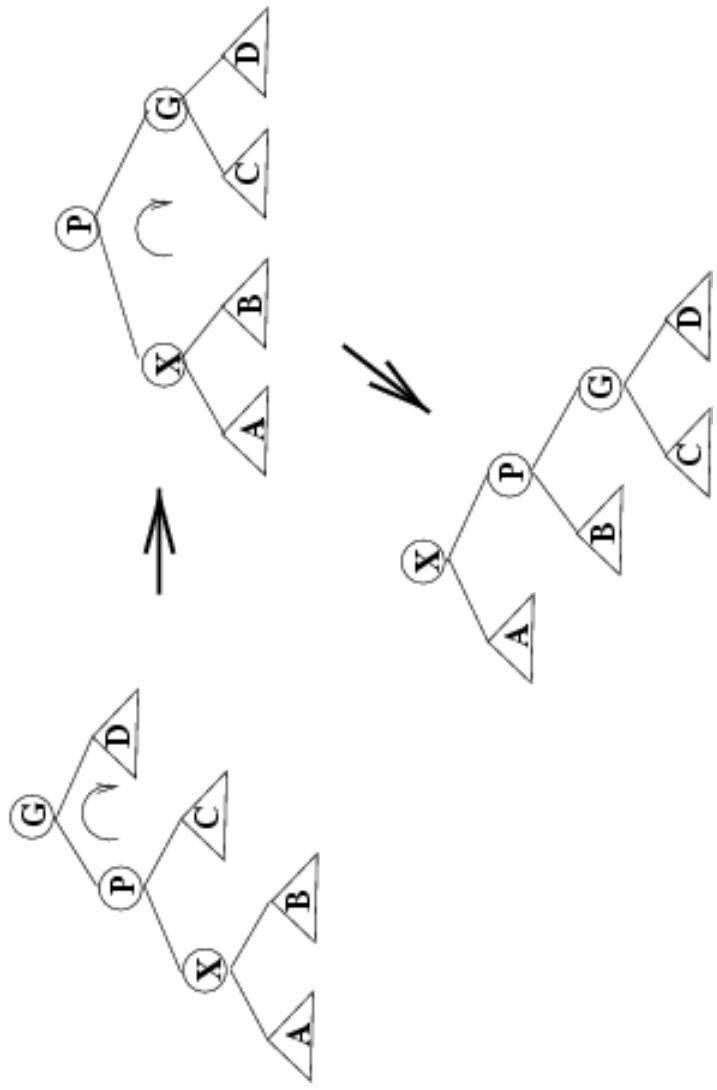
Each rotation

- If  $x$  is root, do nothing.
- If  $x$  has no grandparent, rotate  $x$  about its parent.
- If  $x$  has a grandparent,
  - if  $x$  and its parent are both left children or both right children, rotate the parent about the grandparent, then rotate  $x$  about its parent
  - if  $x$  and its parent are opposite type children (one left and the other right), rotate  $x$  about its parent, then rotate  $x$  about its new parent (former grandparent)

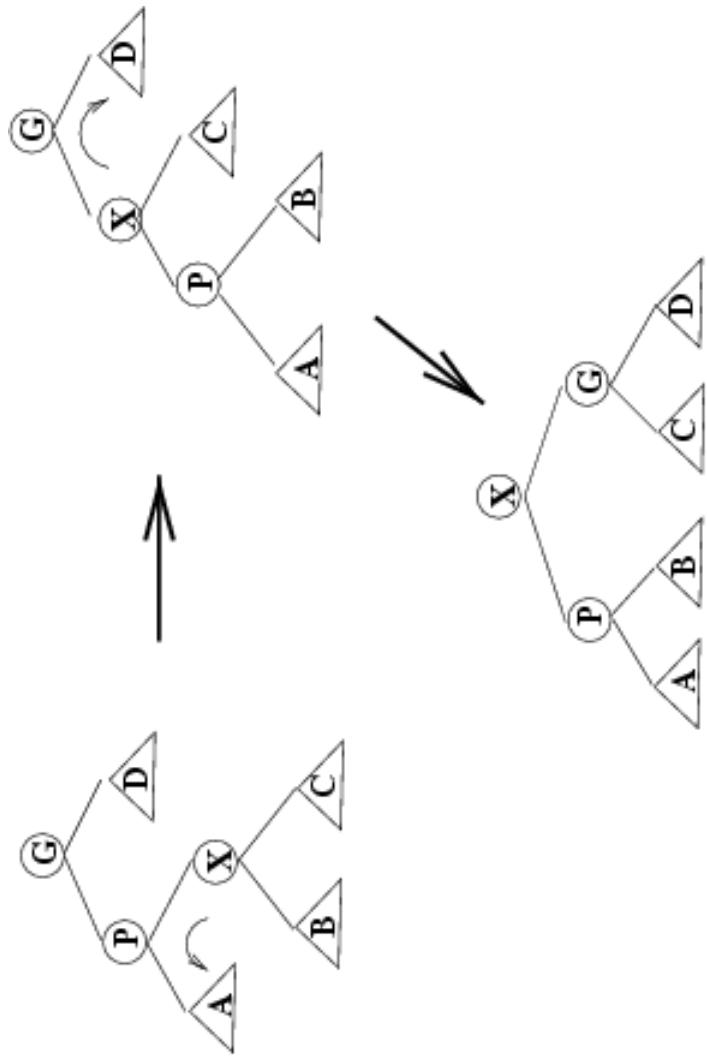
Node has no grandparent



# Node and Parent are Same Side Zig-Zig



# Node and Parent are Different Sides Zig-Zag



# Operations in Splay Trees

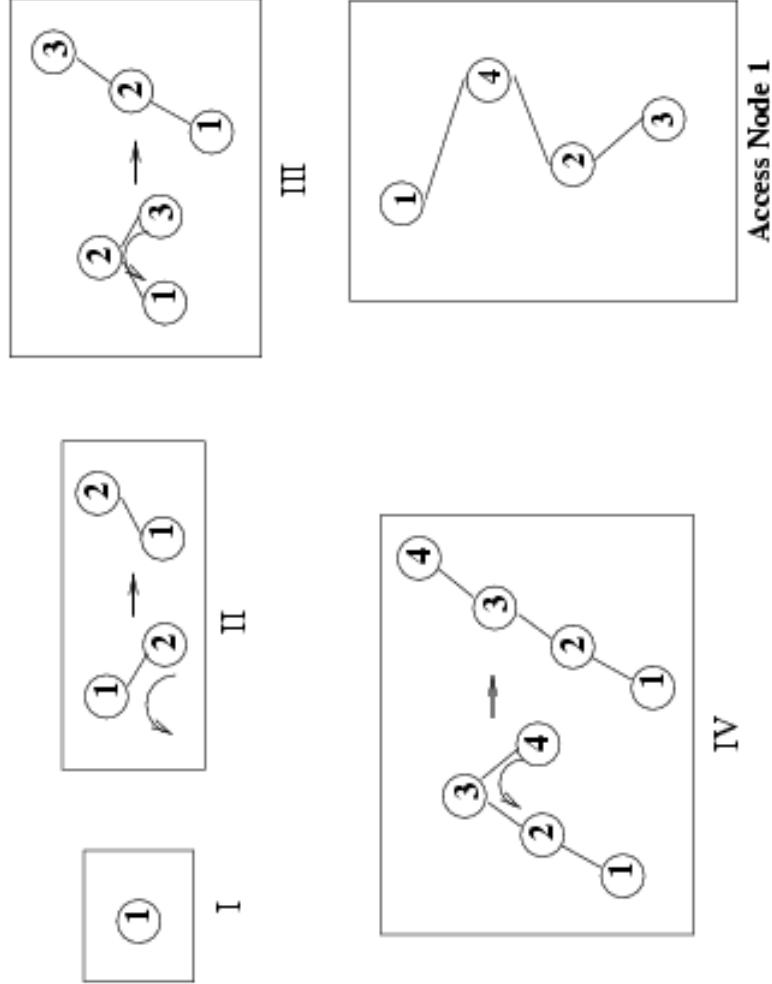
insert

- first insert as in normal binary search tree
- then splay inserted node
- if there is a duplicate, the node holding the duplicate element is splayed

find

- search for node
- if found, splay; otherwise splay last node accessed on the search path

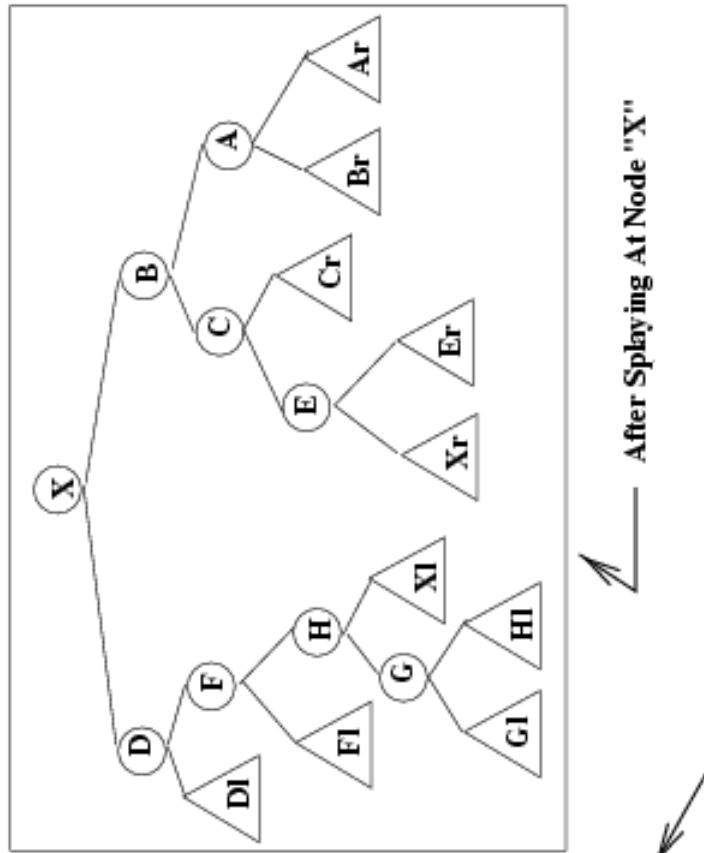
## Insertion in order into a Splay Tree



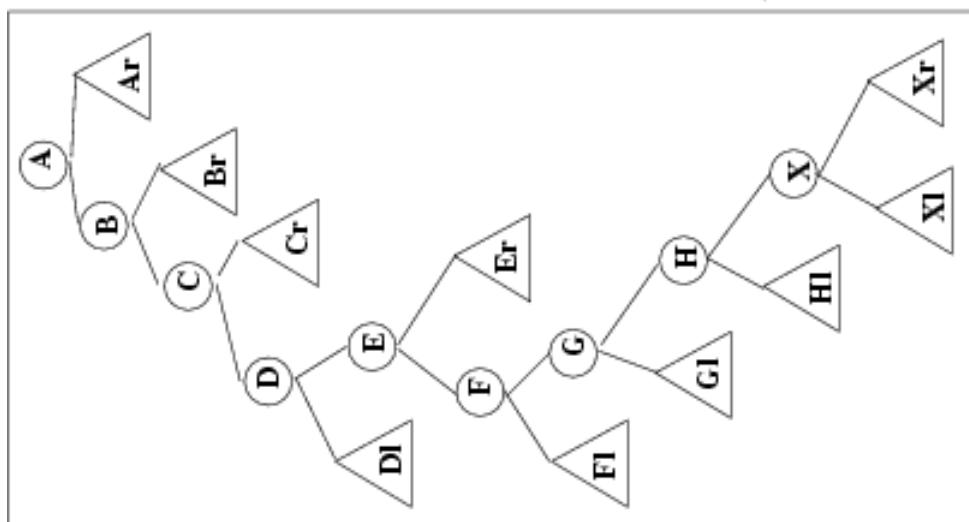
# Operations on Splay Trees (cont)

remove

- splay element to be removed
  - if the element to be deleted is not in the tree, the node last visited on the search path is splayed
- disconnect left and right subtrees from root
- do one of:
  - splay max item in  $T_L$  (then  $T_L$  has no right child)
  - splay min item in  $T_R$  (then  $T_R$  has no left child)
- connect other subtree to empty child of root



Original Tree

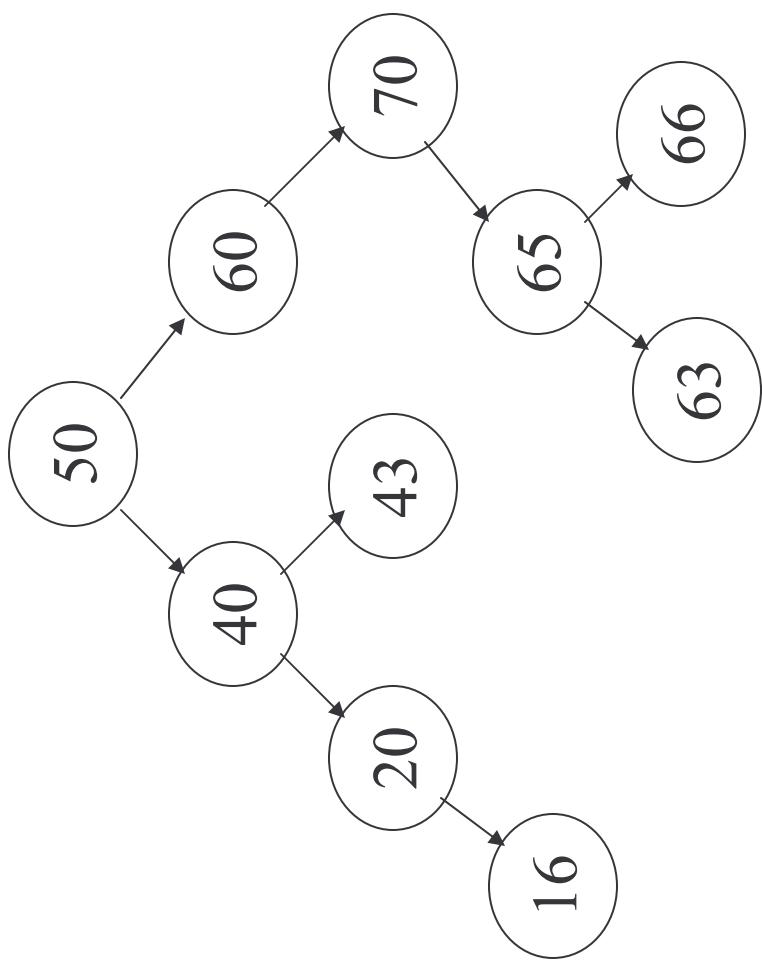


# Performance of Splay Trees

insert

- regular BST insertion --  $O(\text{height})$
- splay:  $O(1)$  for each rotation,  $O(\text{height})$  rotations

Exercise: find( 65 )



Exercise: remove( 25 )

