

Red Black Trees

Top-Down Insertion

Review of Bottom-Up Insertion

- In B-Up insertion, “ordinary” BST insertion was used, followed by correction of the tree on the way back up to the root
- This is most easily done recursively
 - Insert winds up the recursion on the way down the tree to the insertion point
 - Fixing the tree occurs as the recursion unwinds

Top-Down Insertion Strategy

- In T-Down insertion, the corrections are done while traversing down the tree to the insertion point.
- When the actual insertion is done, no further corrections are needed, so no need to traverse back up the tree.
- So, T-Down insertion can be done iteratively which is generally faster

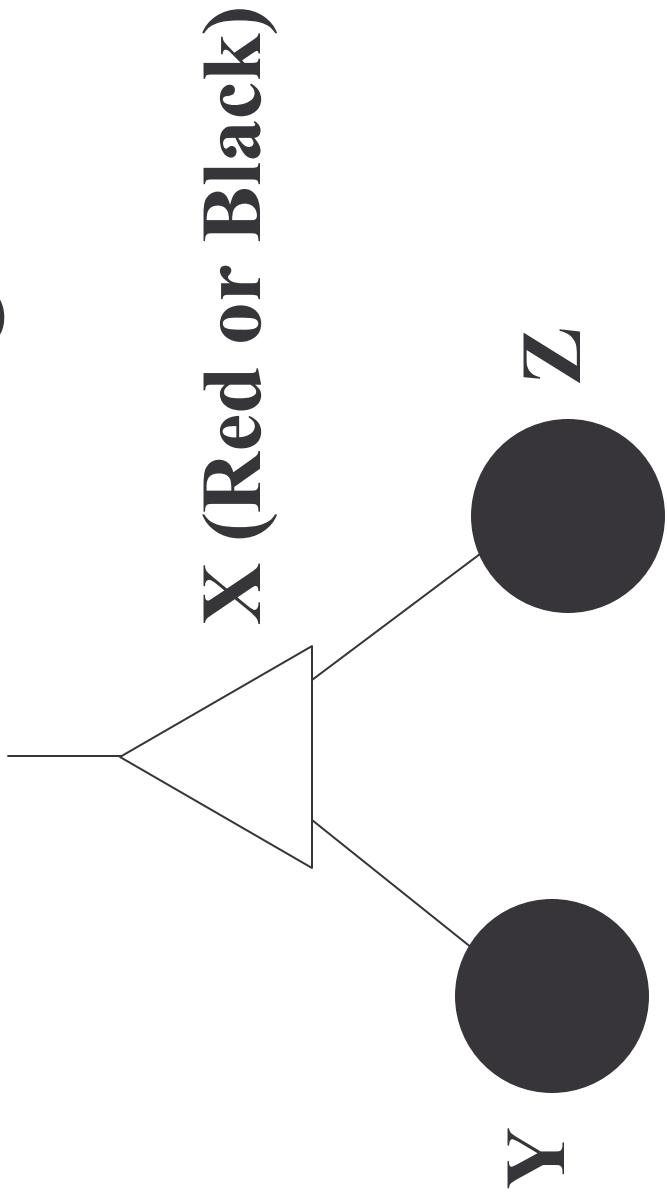
Goal of T-D Insertion

- Insertion is always done as a leaf (as in ordinary BST insertion)
- Recall from the B-Up flow chart that if the uncle of a newly inserted node is Black, we restore the RB tree properties by one or two local rotations and recoloring – we do not need to make changes further up the tree

Goal (2)

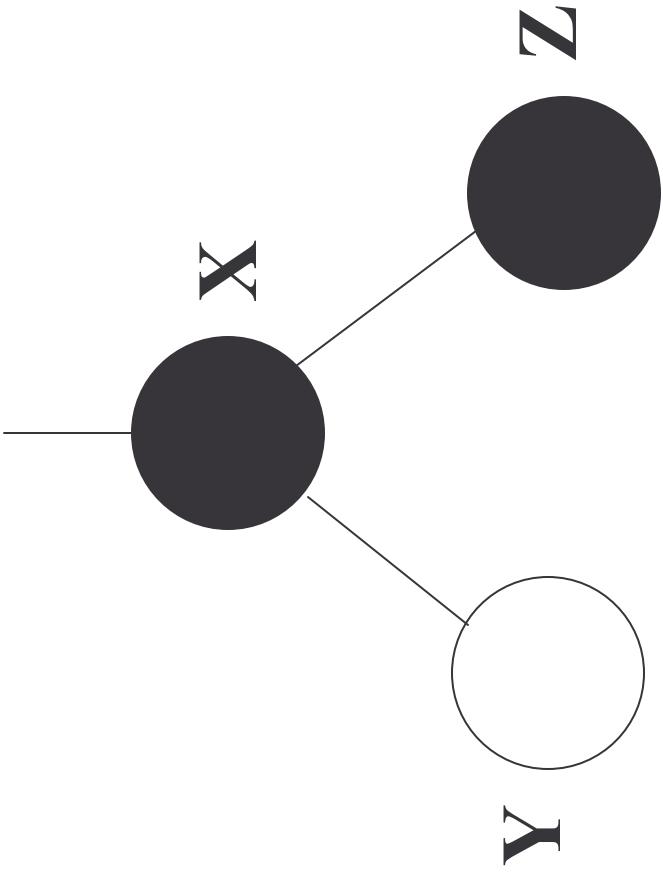
- Therefore, the goal of T-D insertion is to traverse from the root to the insertion point in such a way that RB properties are maintained, and at the insertion point, the uncle is Black.
- That way we may have to rotate and recolor, but not propagate back up the tree

Possible insertion configurations



If a new node is inserted as a child of Y or Z, there is no problem since the new node's parent is Black

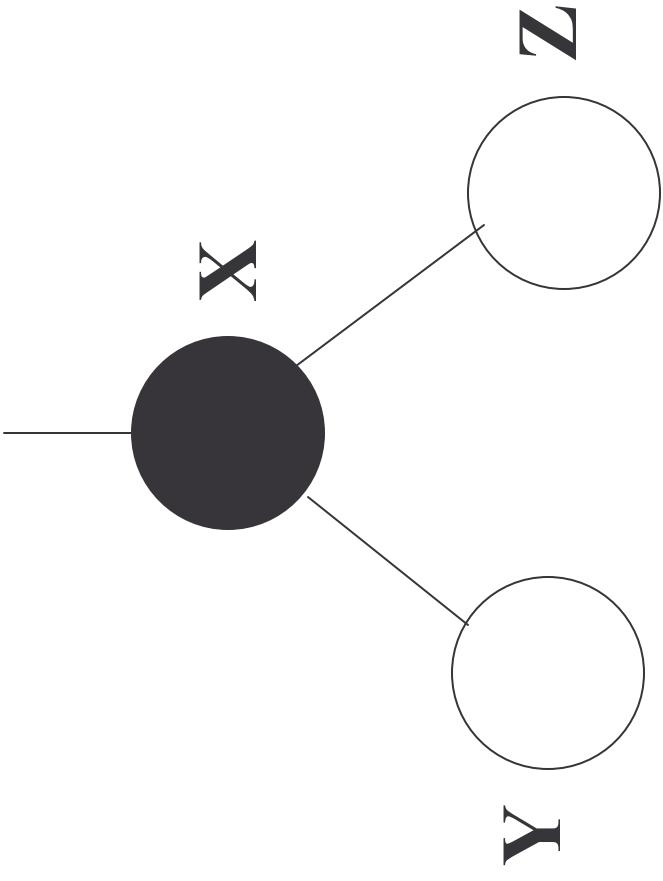
Possible insertion configurations



If new node is child of Z, no problem since Z is Black.

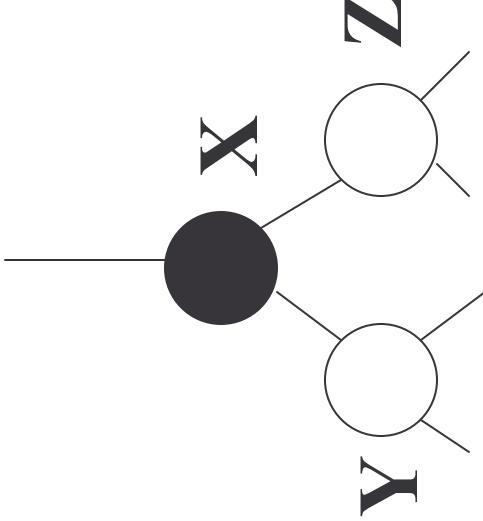
If new node is child of Y, fixable problem since the new node's uncle (Z) is Black – do a few rotations and recolor.... done

Possible insertion configurations



If new node is inserted as child of Y or Z, its uncle will be Red and we will have to go back up the tree. This is the only case we need to avoid.

Top-Down Traversal¹

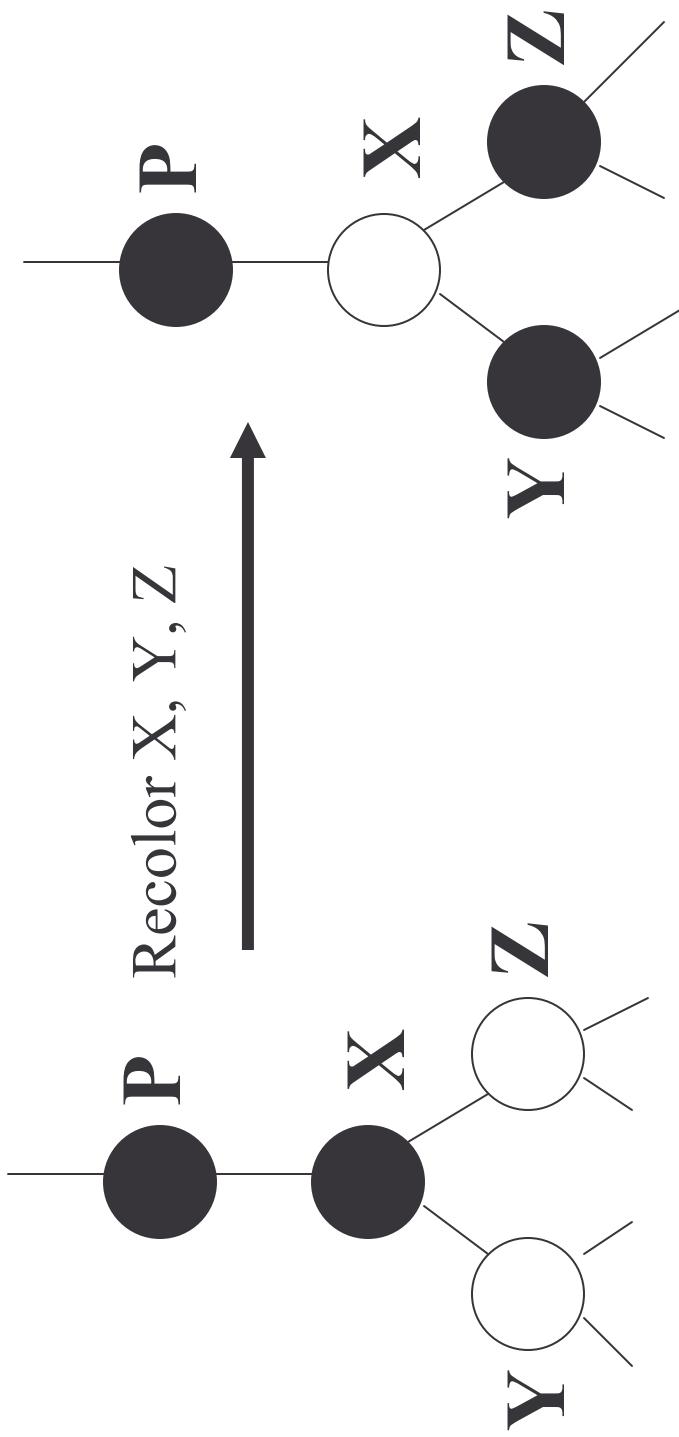


As we traverse down the tree and encounter this case, we recolor and possibly do some rotations.

There are 3 cases.

Remember the goal – to create an insertion point at which the parent of the new node is Black, or the uncle of the new node is Black.

Case 1 – X’s Parent is Black

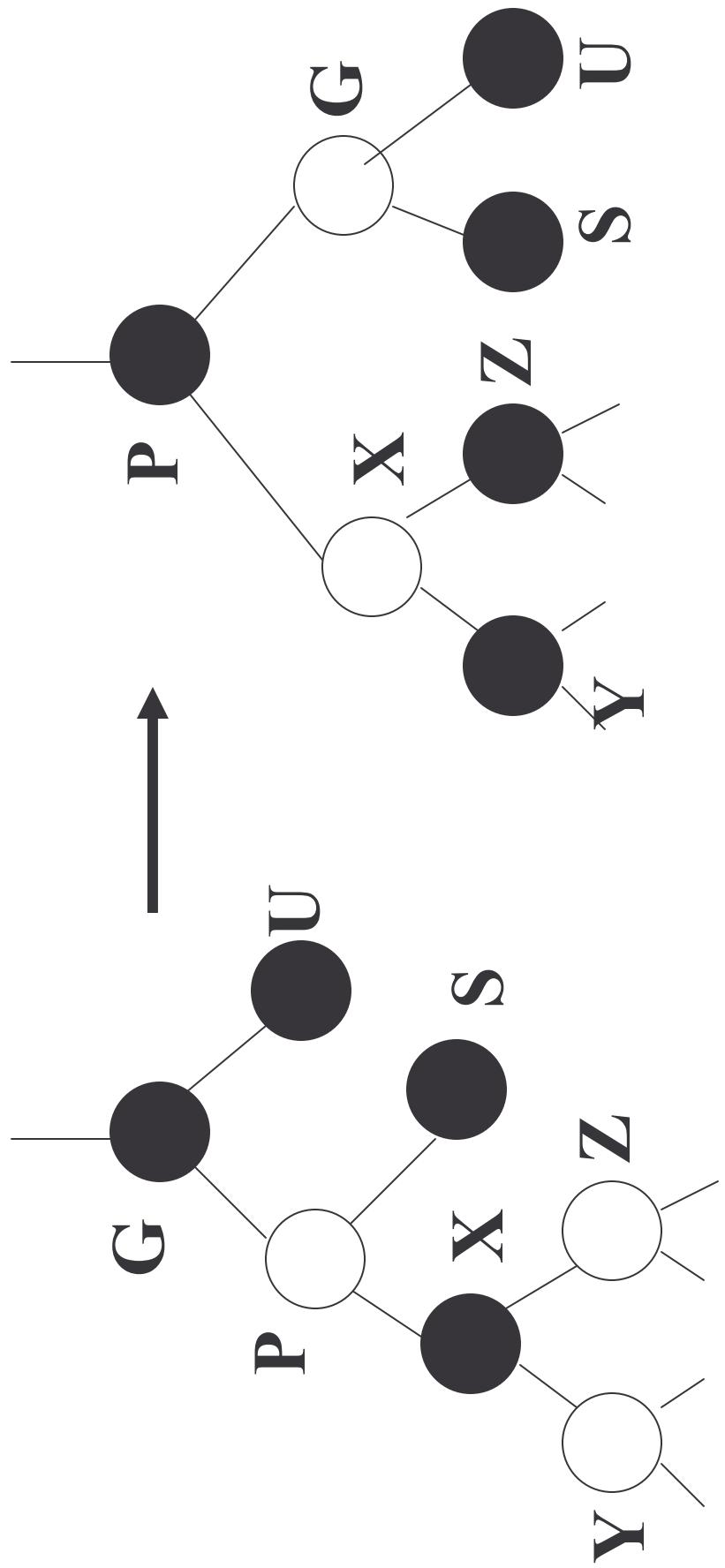


Just recolor and continue down the tree

Case 2

- X's Parent is Red (so Grandparent is Black) and X and P are both left/right children
 - Rotate P around G
 - Color P Black
 - Color G Red
- Note that X's uncle, U, must be Black because it (a) was initially Black, or (b) would have been made Black when we encountered G (which would have had two Red children -- X's Parent and X's uncle)

Case 2 diagrams

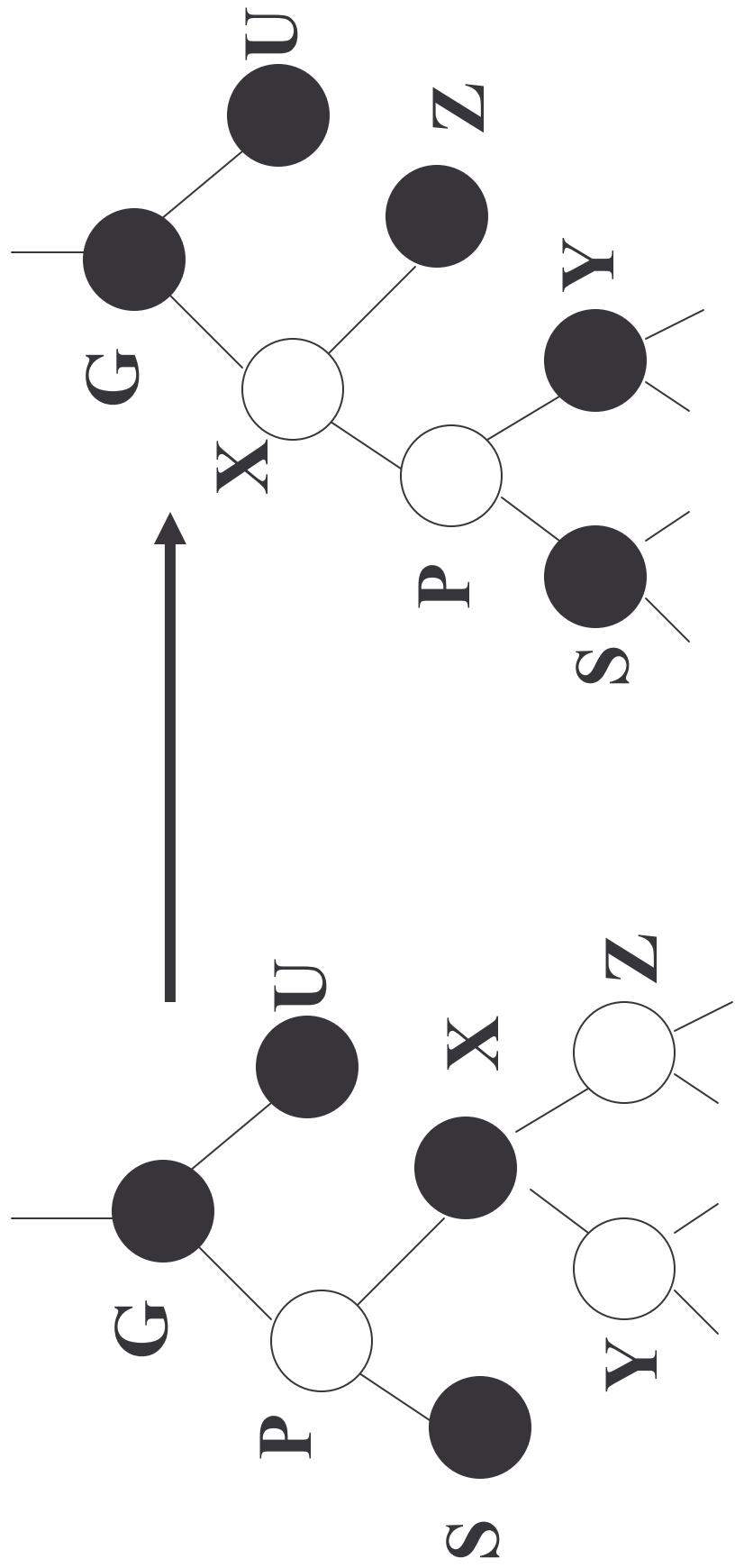


Rotate P around G. Recolor X, Y, Z, P and G

Case 3

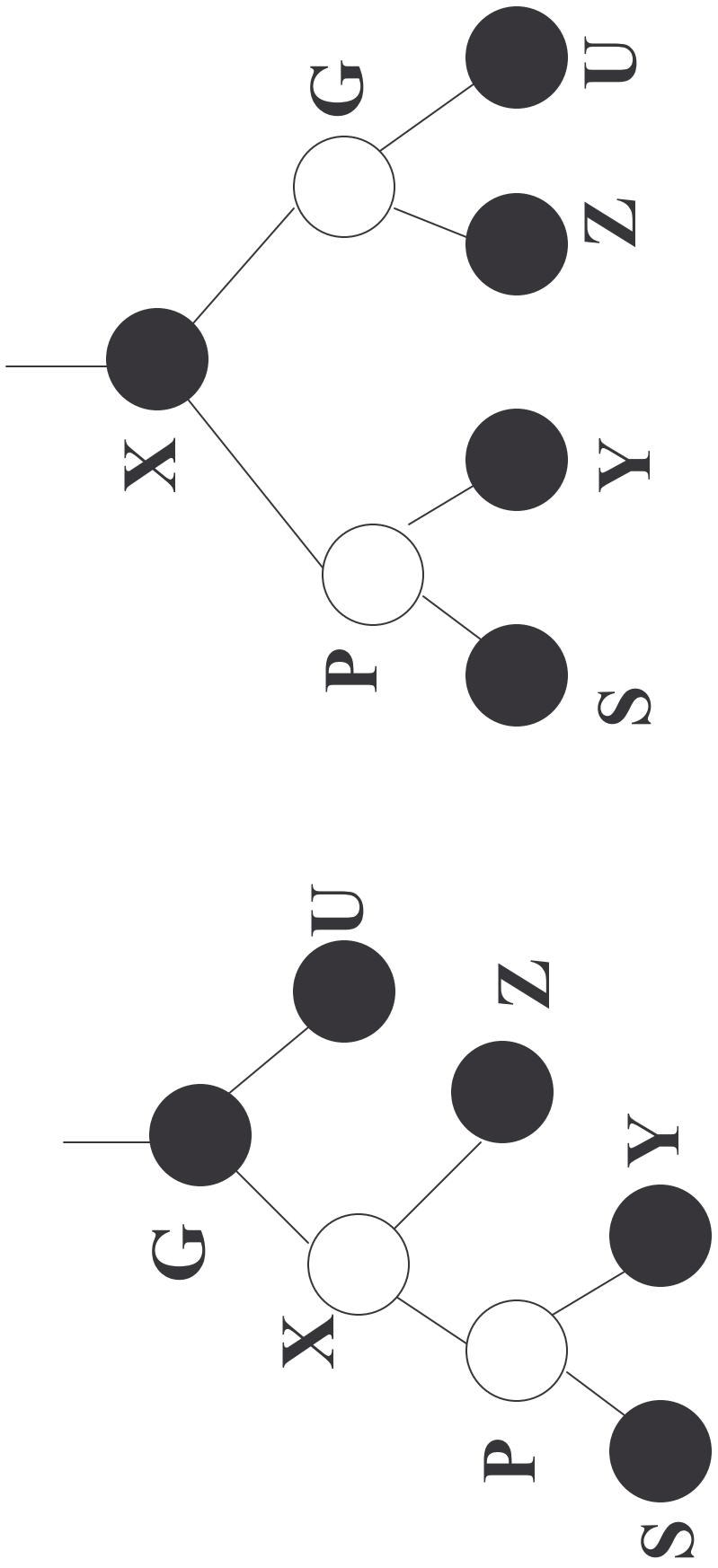
- X's Parent is Red (so Grandparent is Black) and X and P are opposite children
 - Rotate P around G
 - Color P Black
 - Color G Red
- Again note that X's uncle, U, must be Black because it (a) was initially Black, or (b) would have been made Black when we encountered G (which would have had two Red children -- X's Parent and X's uncle)

Case 3 Diagrams (1 of 2)



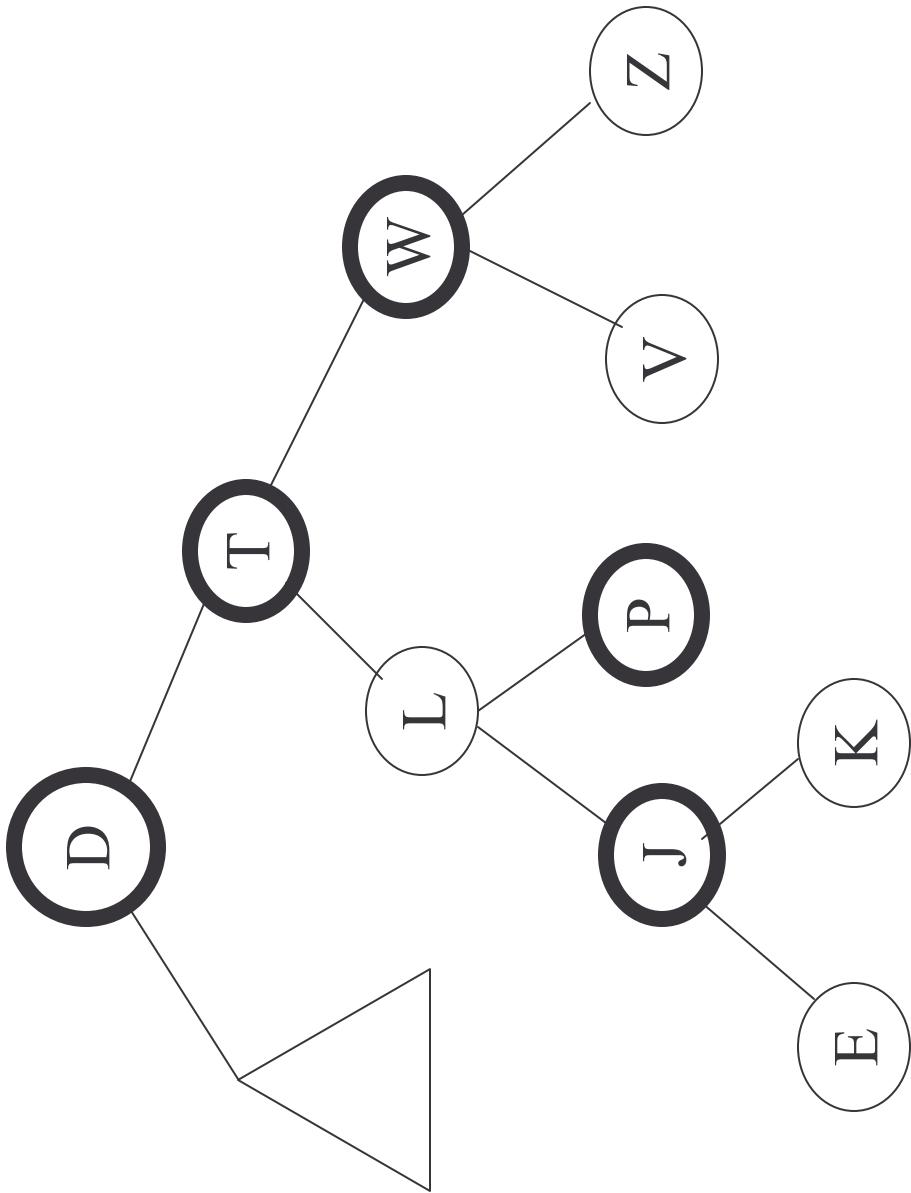
Step 1 – recolor X, Y and Z. Rotate X around P.

Case 3 Diagrams (2 of 2)



Step 2 – Rotate X around G. Recolor X and G

An exercise – insert F

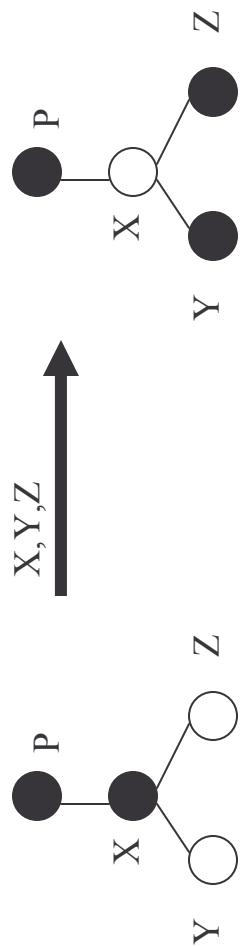


Top-Down Insert Summary

Case 1

Pis Black

Just Recolor



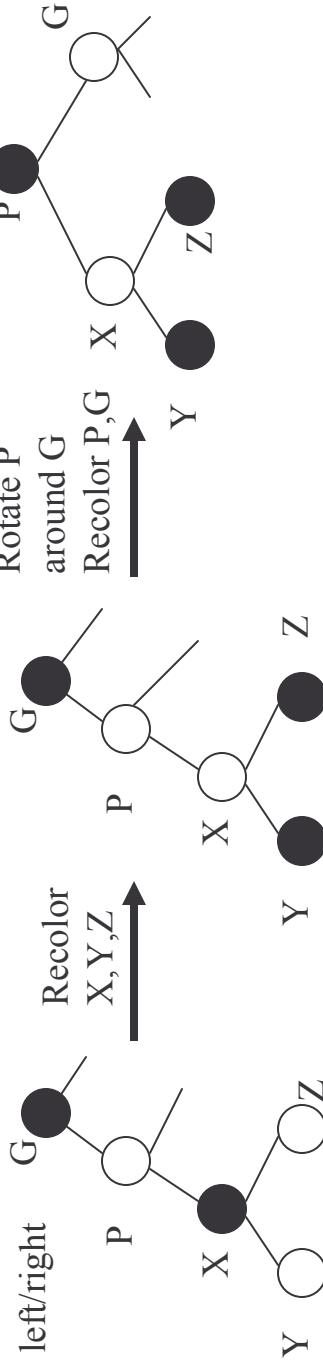
Recolor
X,Y,Z

P
X,Y,Z

Case 2

P is Red

X & P both left/right



Dado

Bassarow

G

Rotate P around G

5

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P is Red

X and P are

and are
omosite chi

Recolor X G

10 of 10

Z
Y

Z
Y

Z